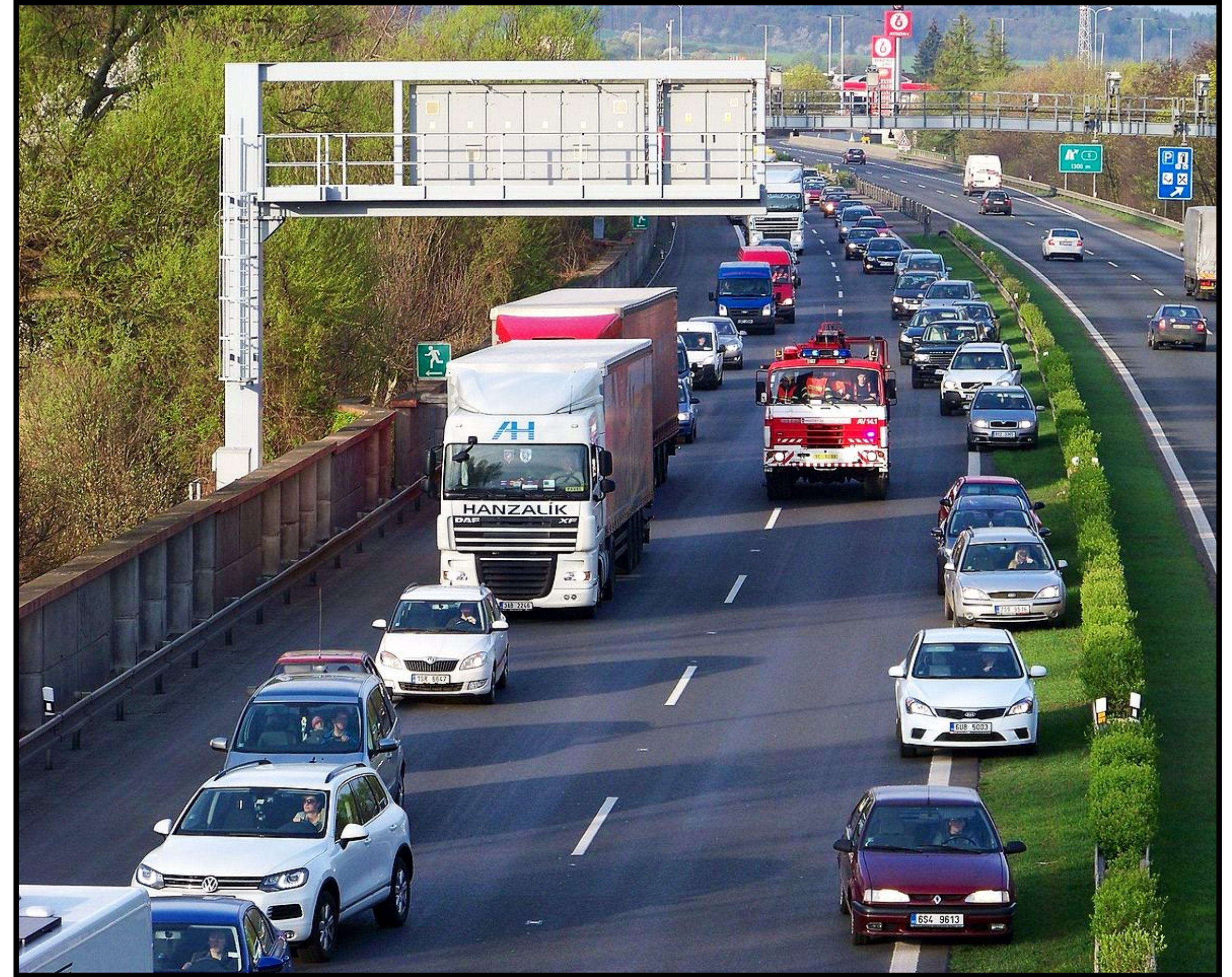




In-Network Congestion Management for Security and Performance

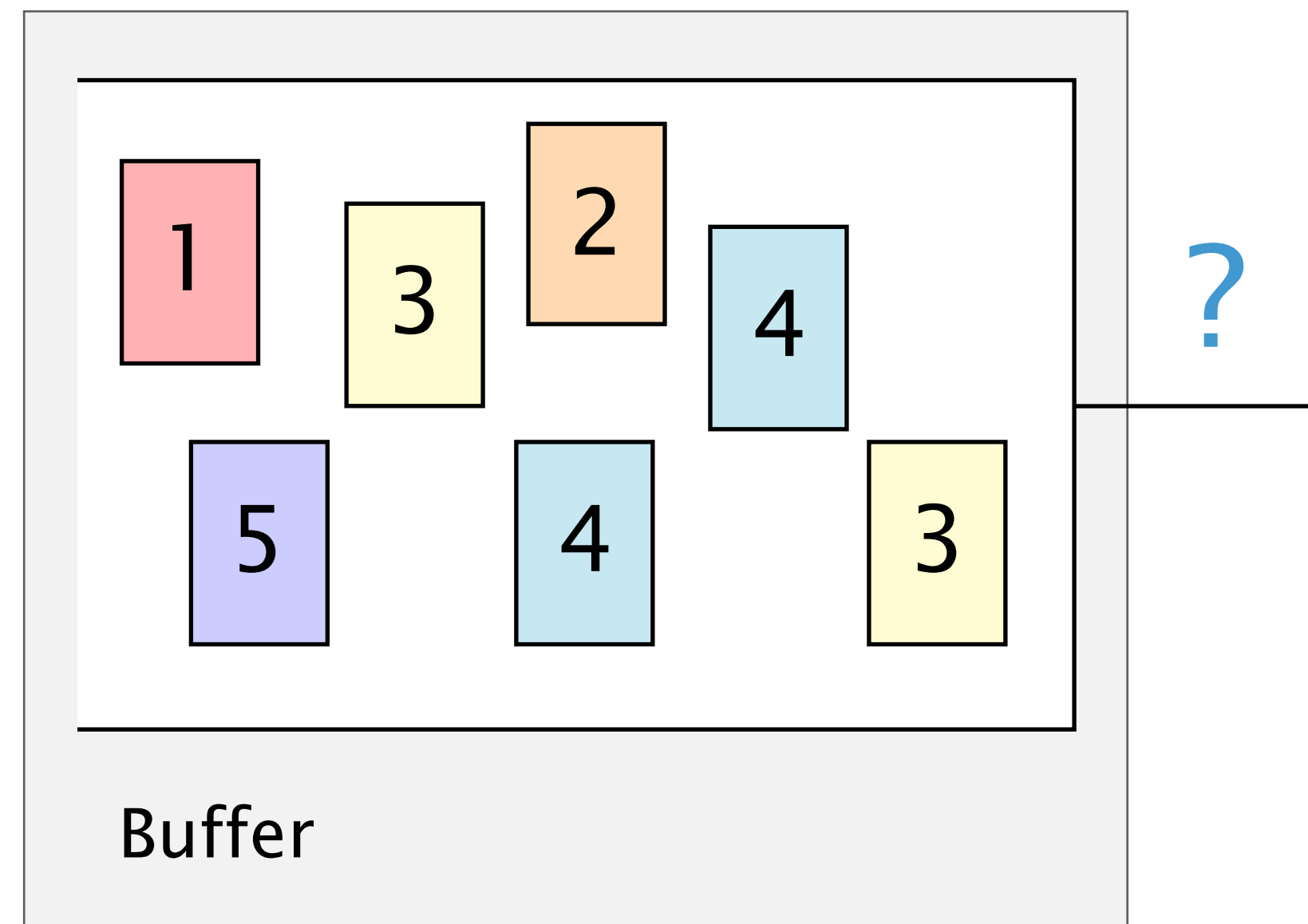
Albert Gran Alcoz
September 18 2024





Packet scheduling

Packet scheduling defines **what packet** should we send next and **when**



Researchers have proposed dozens of scheduling algorithms

Minimize flow completion times

Prioritize packets from short flows

SRPT, PIAS

Enforce fairness

Send one packet from each class at a time

RR, WFQ

Minimize tail latency

Prioritize packets with high slack time

FIFO+, LSTF

A universal scheduling algorithm does **not** exist

NSDI'16

Universal Packet Scheduling

Radhika Mittal[†] Rachit Agarwal[†] Sylvia Ratnasamy[†] Scott Shenker^{†‡}
[†]UC Berkeley [‡]ICSI

Abstract

In this paper we address a seemingly simple question: *Is there a universal packet scheduling algorithm?* More precisely, we analyze (both theoretically and empirically) whether there is a single packet scheduling algorithm that, at a network-wide level, can perfectly match the results of *any* given scheduling algorithm. We find that in general the answer is “no”. However, we show theoretically that the classical Least Slack Time First (LSTF) scheduling algorithm comes closest to being universal and demonstrate empirically that LSTF can closely replay a wide range of scheduling algorithms in realistic network settings. We then evaluate whether LSTF can be used *in practice* to meet various network-wide objectives by looking at popular performance metrics (such as mean FCT, tail packet delays, and fairness); we find that LSTF performs comparable to the state-of-the-art for each of them. We also discuss how LSTF can be used in conjunction with active queue management schemes (such as CoDel) without changing the core of the network.

1 Introduction

There is a large and active research literature on novel packet scheduling algorithms, from simple schemes such as priority scheduling [31], to more complicated mech-

We can define a universal packet scheduling algorithm (hereafter UPS) in two ways, depending on our viewpoint on the problem. From a theoretical perspective, we call a packet scheduling algorithm *universal* if it can replay any *schedule* (the set of times at which packets arrive to and exit from the network) produced by any other scheduling algorithm. This is not of practical interest, since such schedules are not typically known in advance, but it offers a theoretically rigorous definition of universality that (as we shall see) helps illuminate its fundamental limits (i.e., which scheduling algorithms have the flexibility to serve as a UPS, and why).

From a more practical perspective, we say a packet scheduling algorithm is universal if it can achieve different desired performance objectives (such as fairness, reducing tail latency, minimizing flow completion times). In particular, we require that the UPS should match the performance of the best known scheduling algorithm for a given performance objective.¹

The notion of universality for packet scheduling might seem esoteric, but we think it helps clarify some basic questions. If there exists no UPS then we should *expect* to design new scheduling algorithms as performance objectives evolve. Moreover, this would make a strong argument for switches being equipped with programmable

How to deploy **all** scheduling algorithms?

✗ Generality

Universal packet scheduler

Flexibility

Customized algorithms

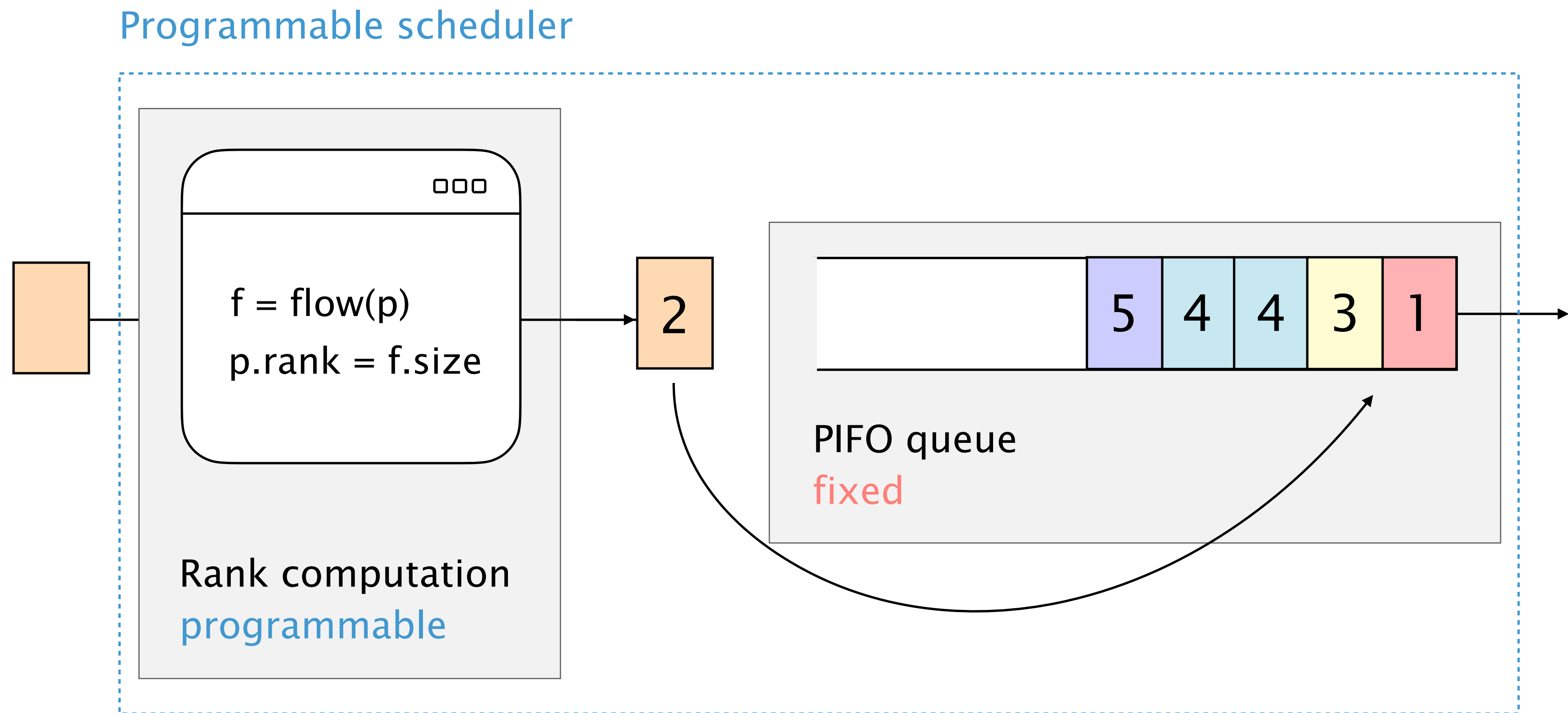
How to deploy **all** scheduling algorithms?

✗ Generality

Universal packet scheduler

Programmable
scheduling

Push-In First-Out (PIFO) queues enable programmable scheduling



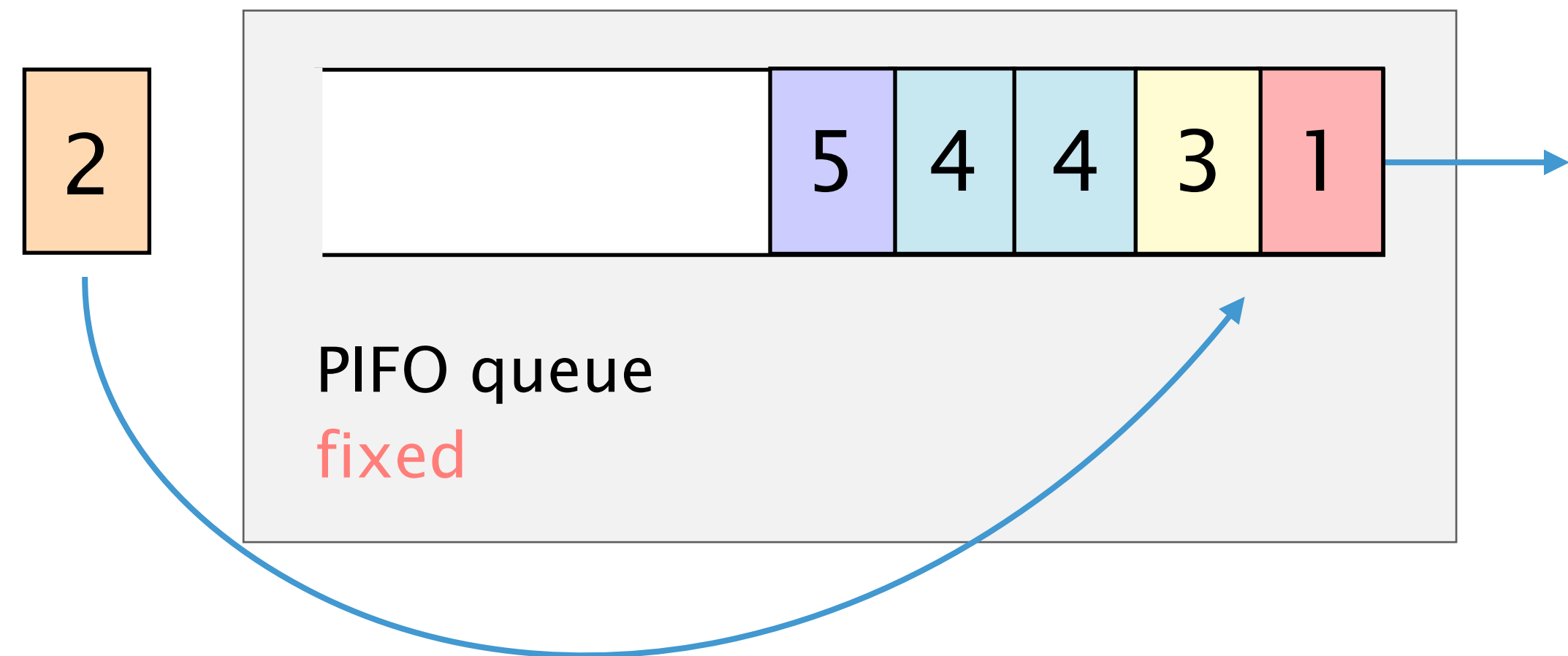
PIFO queues are characterized by **two** key behaviors

Admission

Enqueue packets with the lowest ranks

Scheduling

Forward packets in rank order



How to implement PIFO queues on hardware?

New ASIC

High performance ✓

~200M \$ ✗

Multiple years ✗

How to implement PIFO queues on hardware?

New ASIC

High performance

~200M \$

Multiple years



Programmable switches

~10K \$

Available today



How to implement PIFO queues on hardware?

New ASIC

High performance



~200M \$



Multiple years



Programmable switches

Enough performance



~10K \$



Available today



Objective

Enable programmable scheduling on existing devices
to improve the Internet's performance and security

How to enable programmable scheduling
on existing devices?

SP-PIFO

[NSDI'20]

Approximating
PIFO's scheduling

PACKS

[NSDI'25]

Incorporating
PIFO's admission

How to use it to improve
the Internet's security?

ACC-Turbo

[SIGCOMM'22]

Mitigating
DDoS attacks

How to enable programmable scheduling on existing devices?

SP-PIFO
[NSDI'20]

Approximating
PIFO's scheduling

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[NSDI'25]

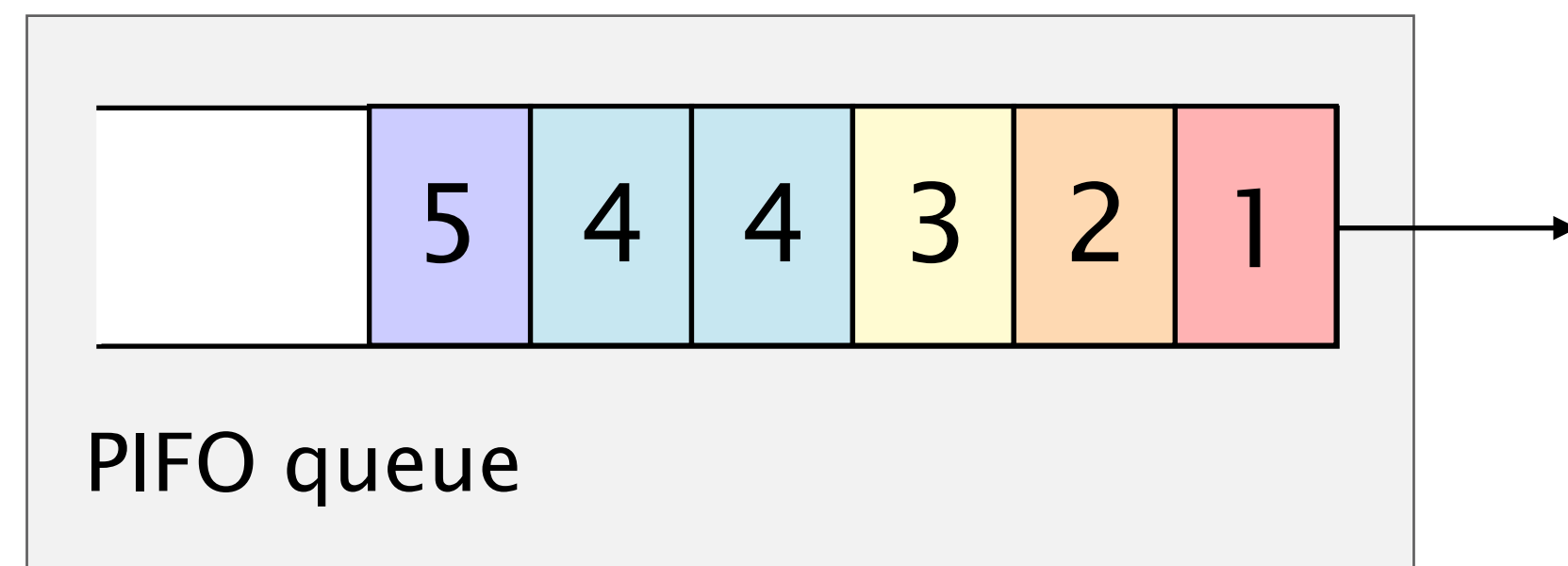
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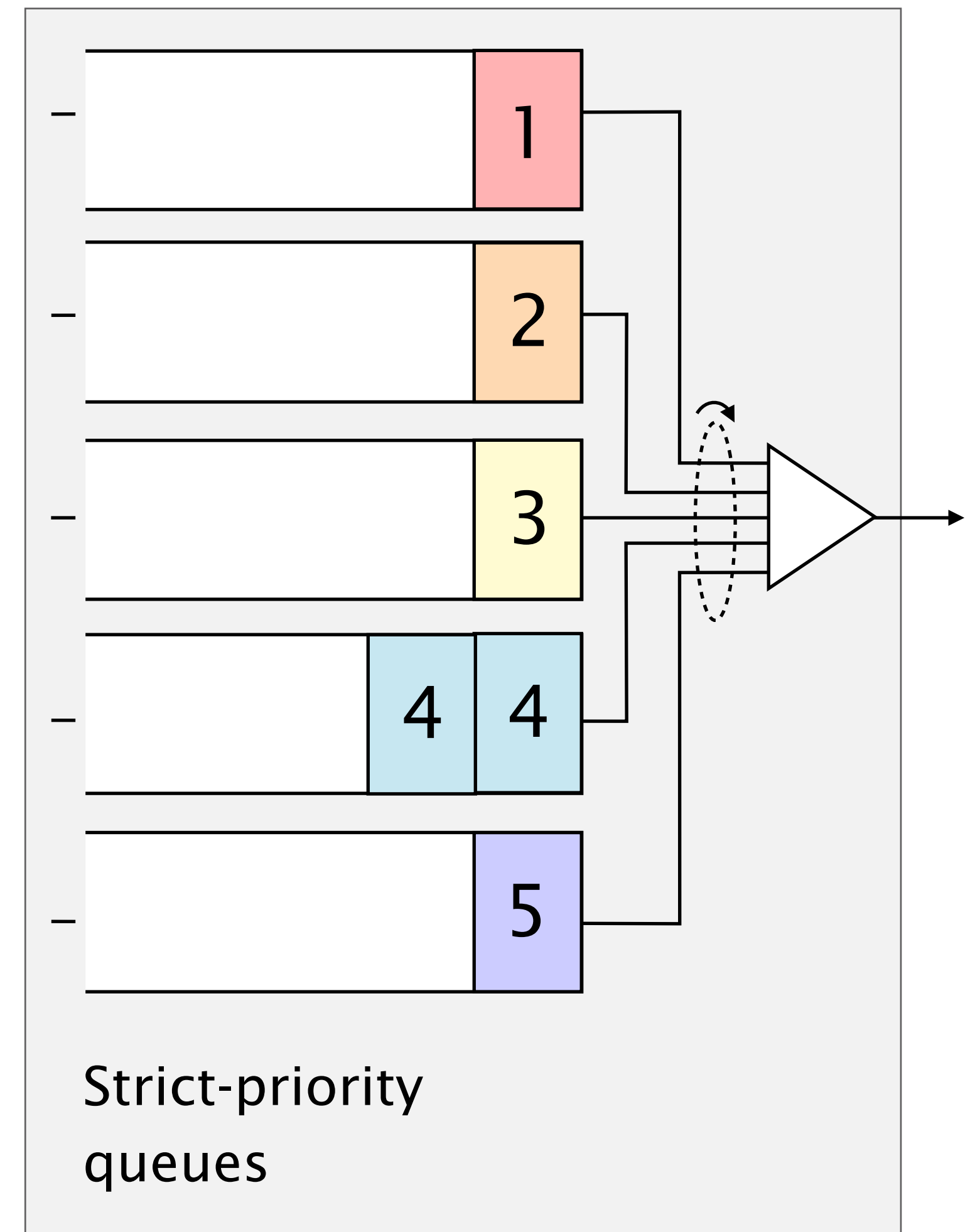
Mitigating
DDoS attacks

We can approximate PIFO queues using strict-priority queues

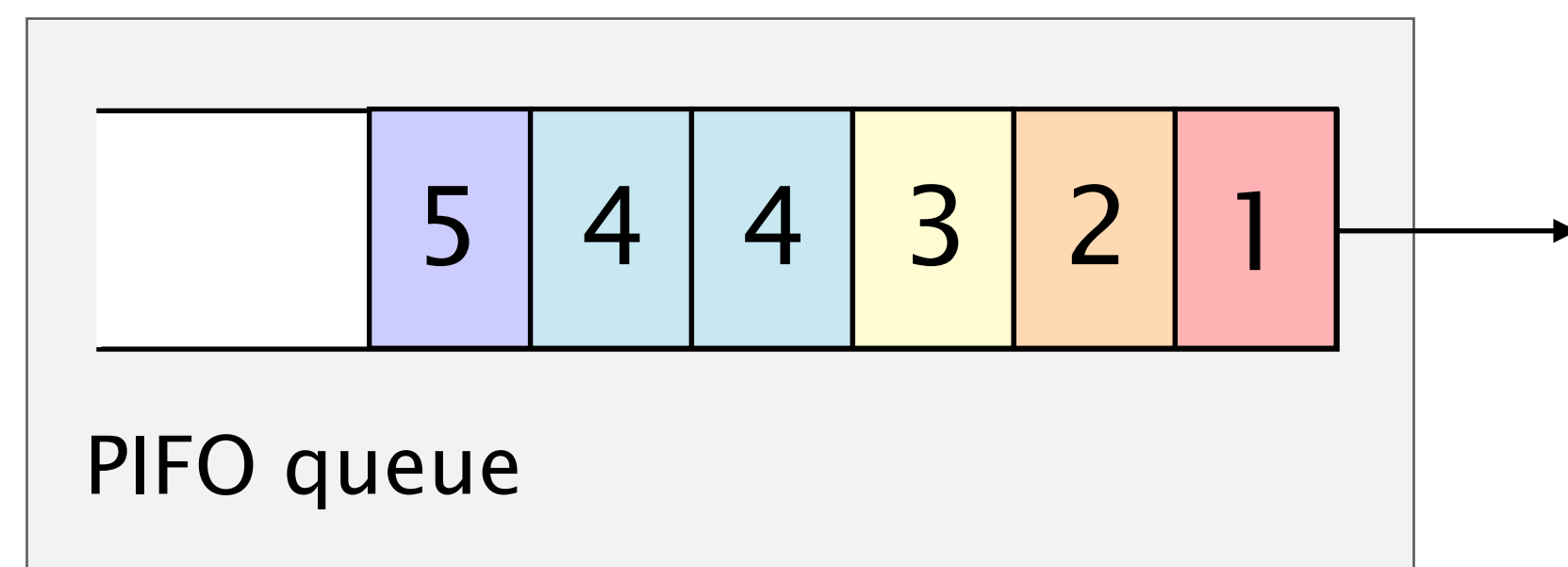


≈

Ideal case One rank per queue



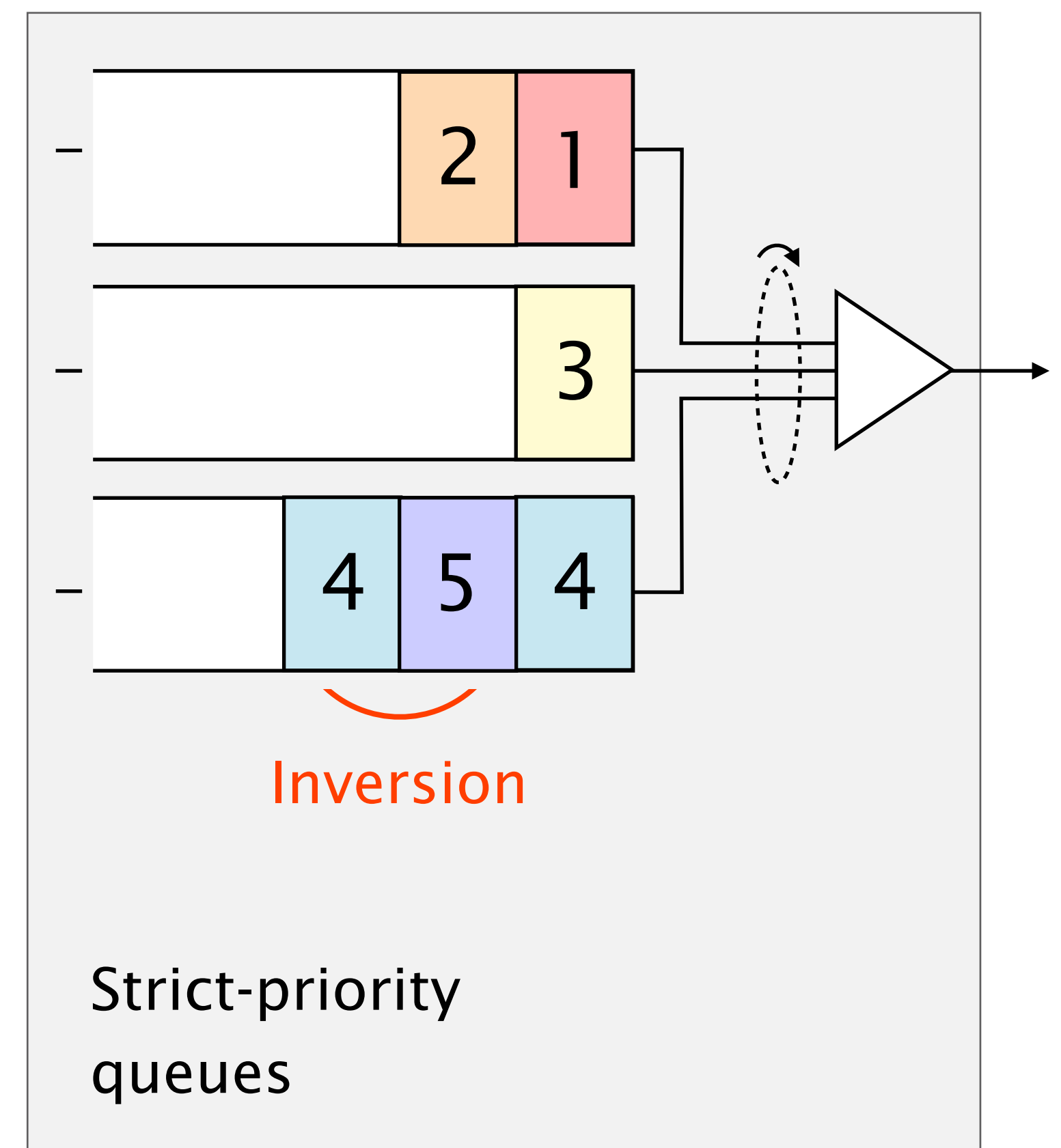
We can approximate PIFO queues using strict-priority queues



≈

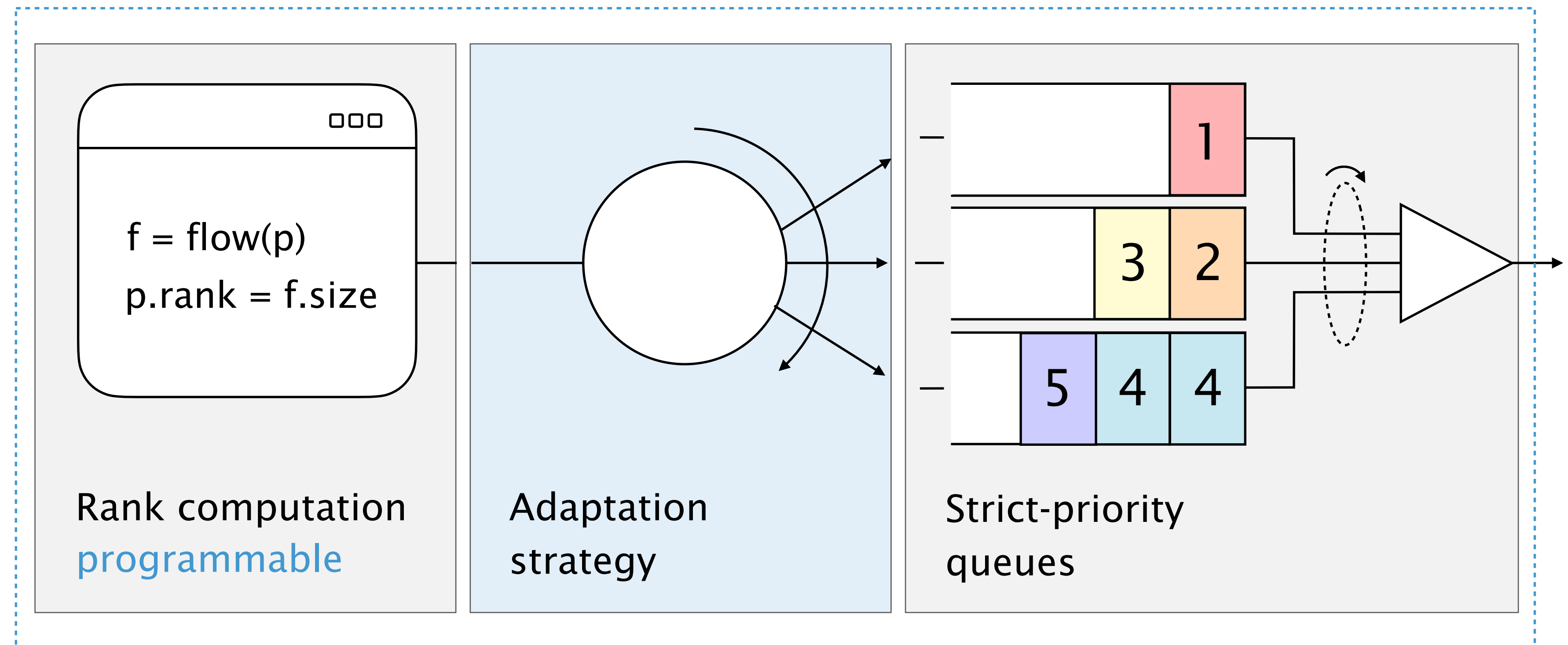
In practice

Multiple ranks per queue



SP-PIFO approximates PIFO queues using strict-priority queues and a dynamic mapping strategy

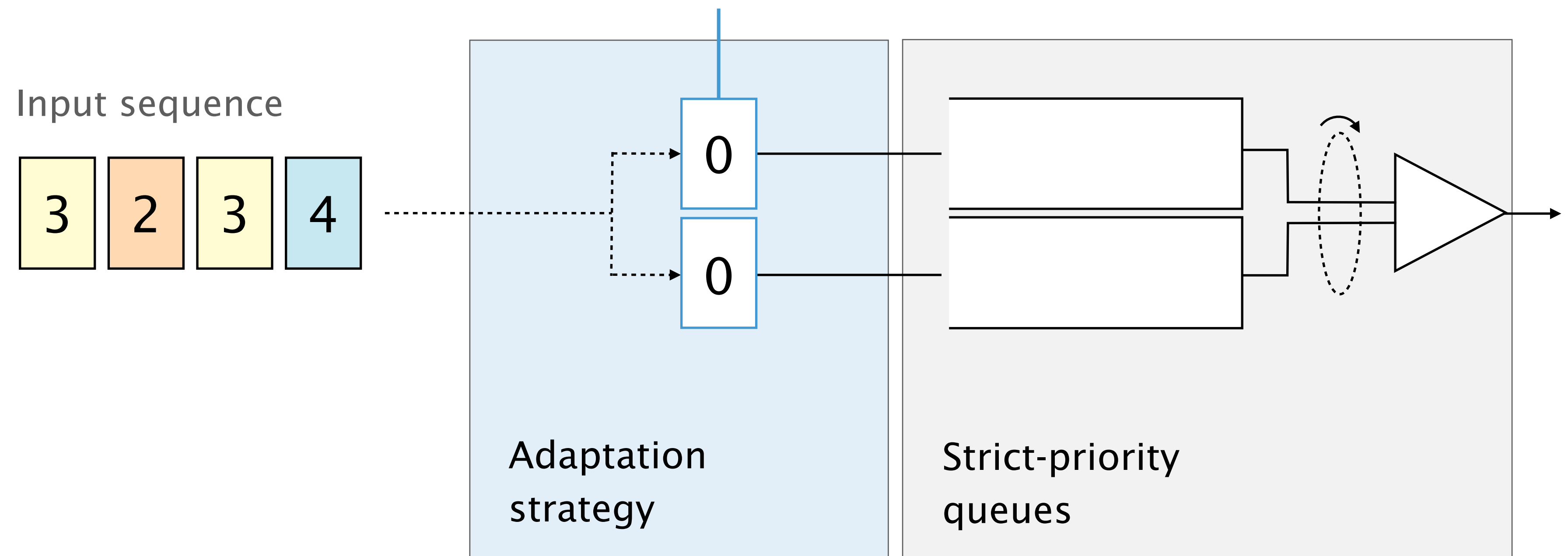
Programmable scheduler



SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Queue bounds

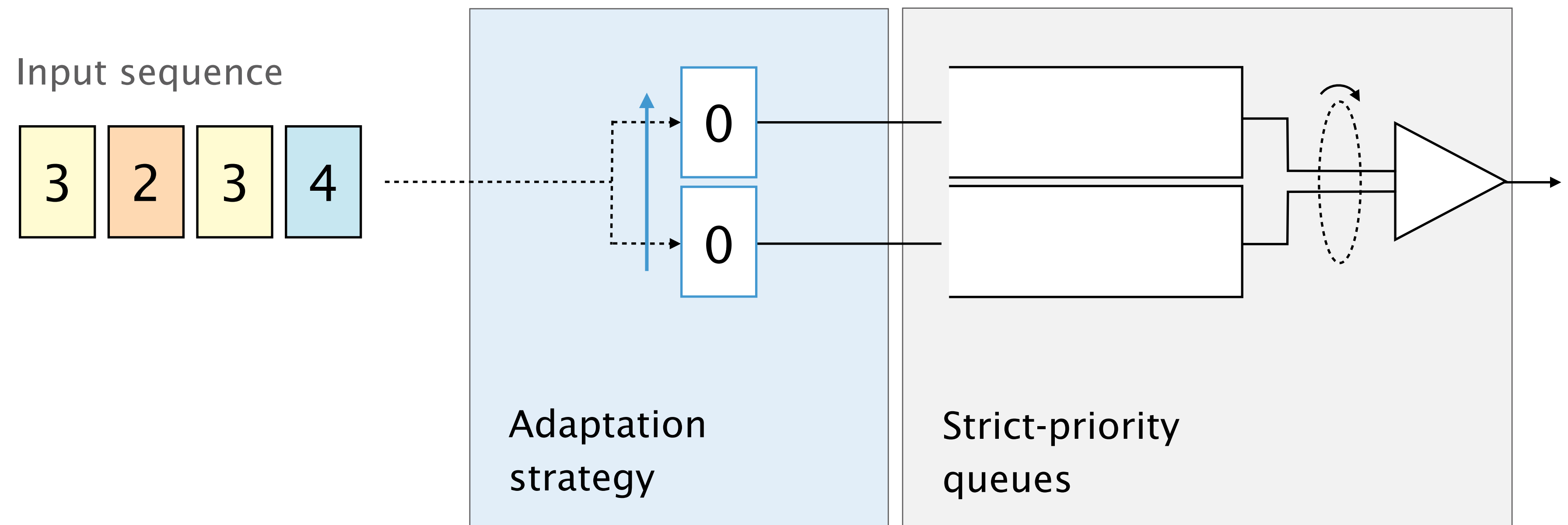
Define which packets to admit to each queue



SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$

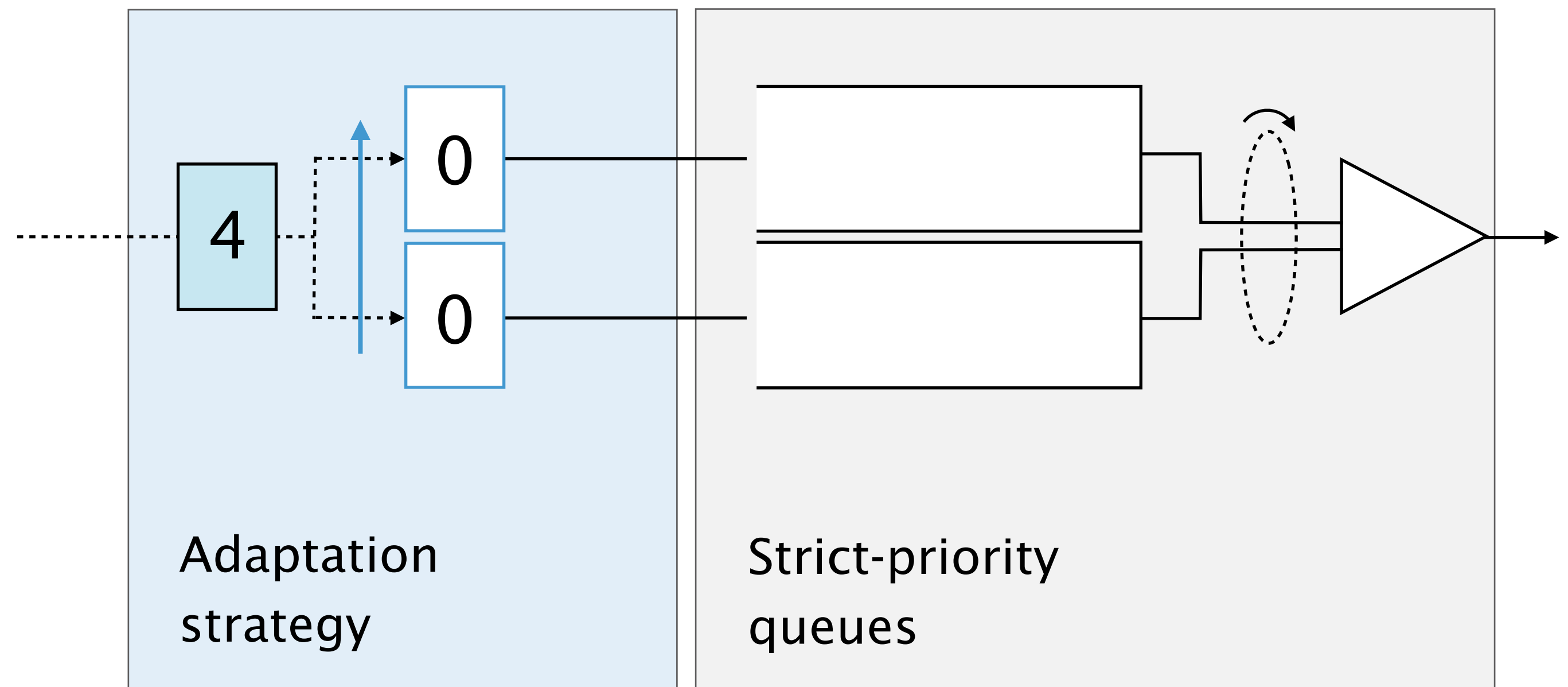
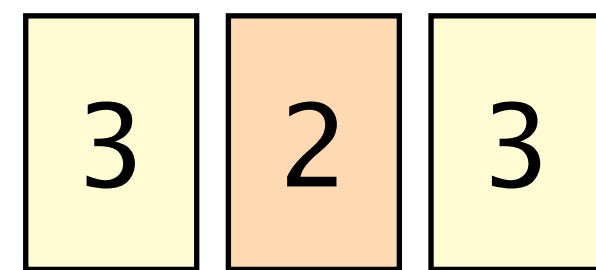


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$

Input sequence

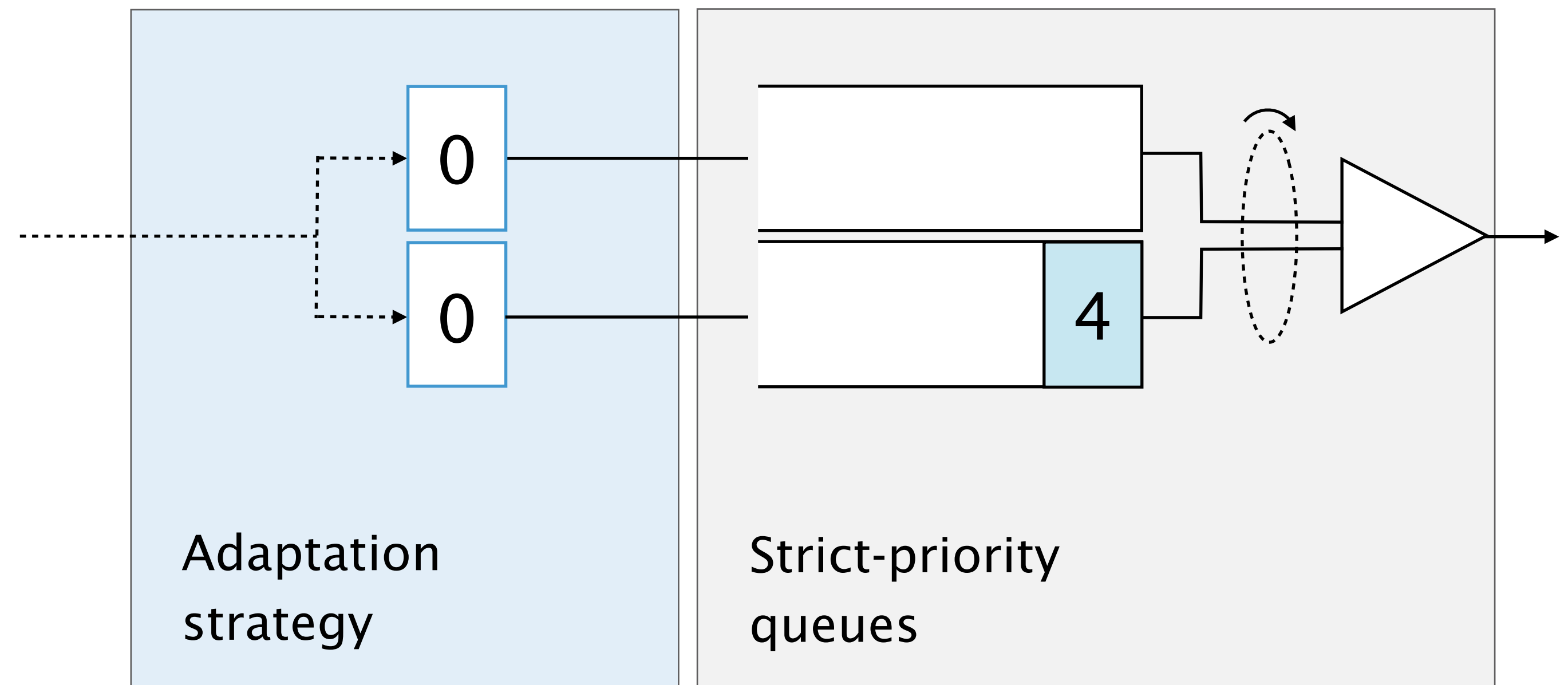
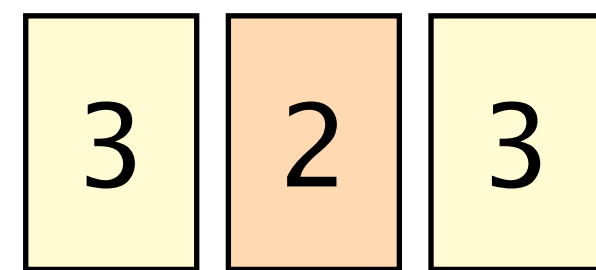


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$

Input sequence

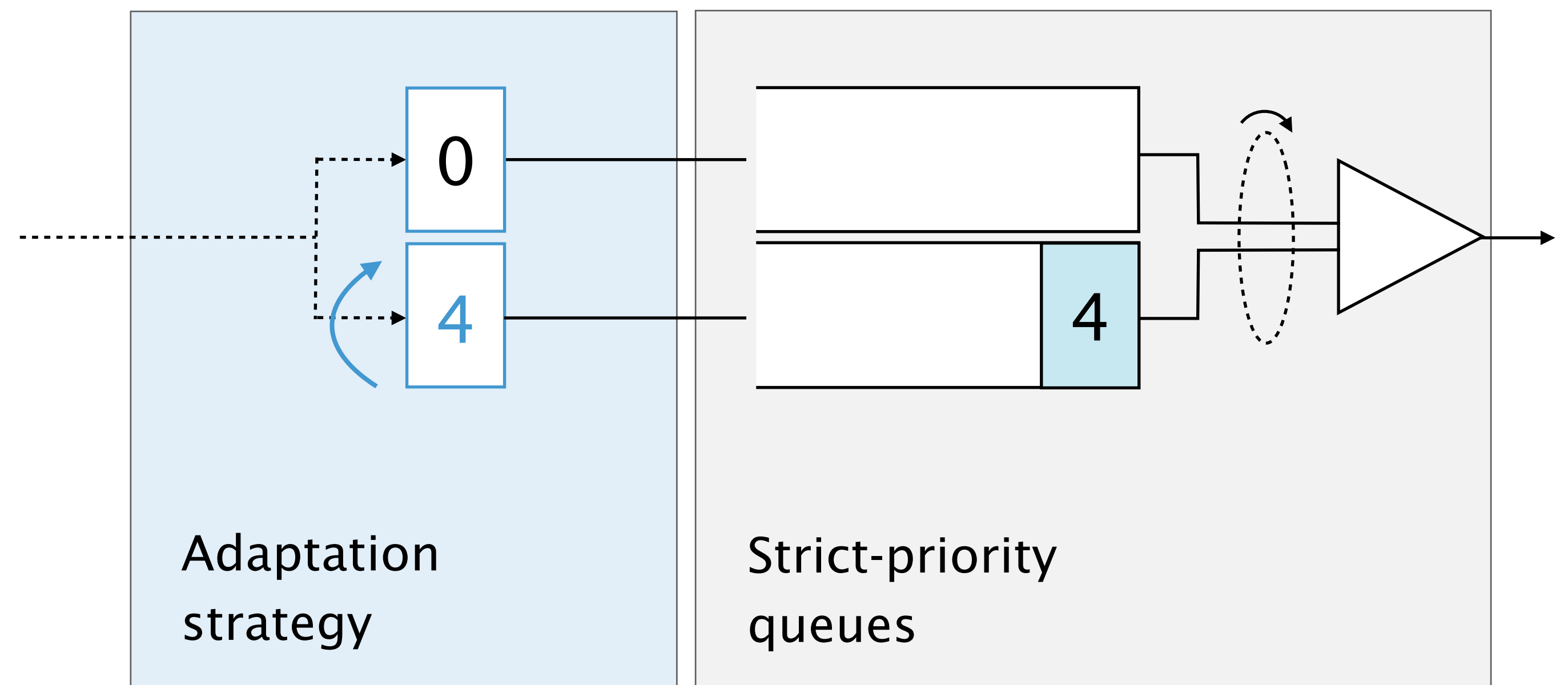
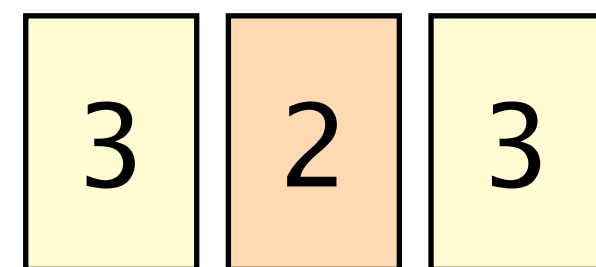


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Push-up adaptation

Set bound to packet rank after enqueue

Input sequence

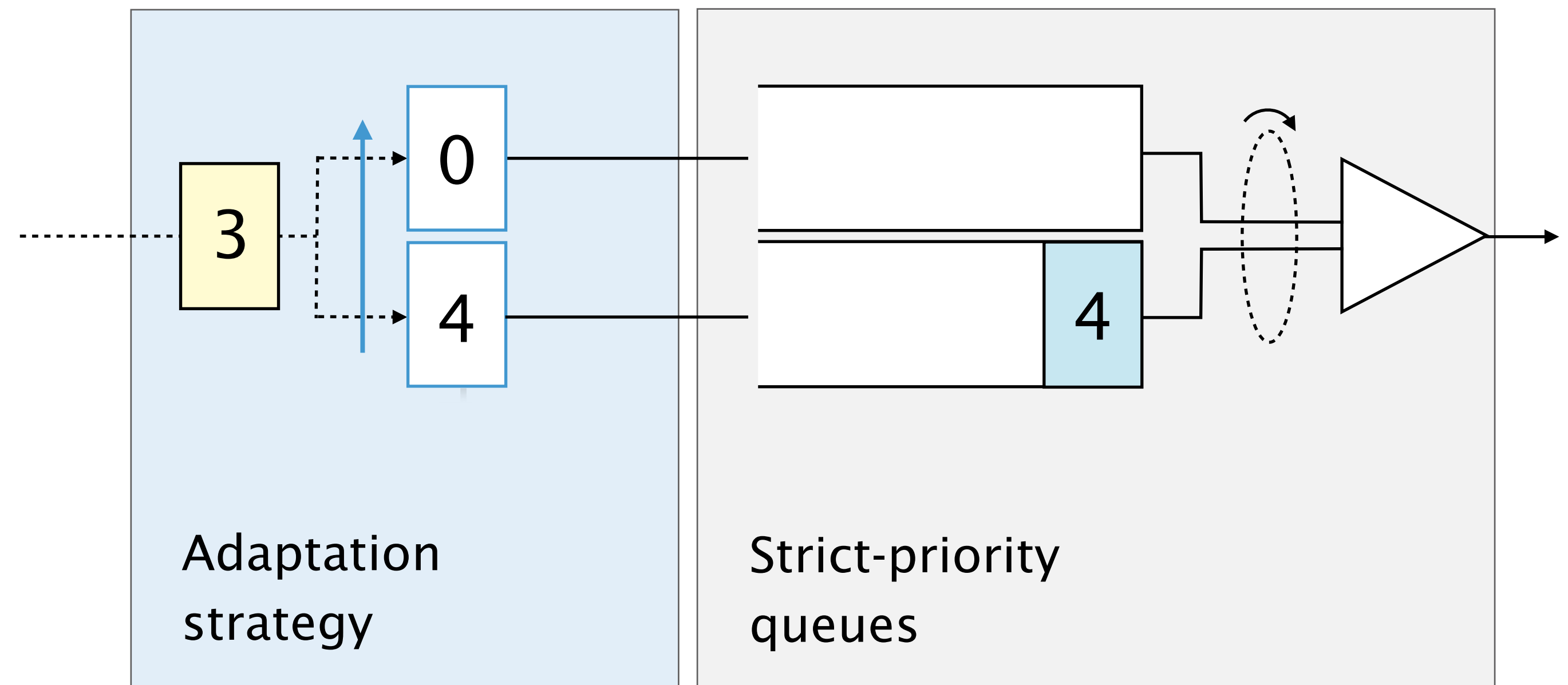
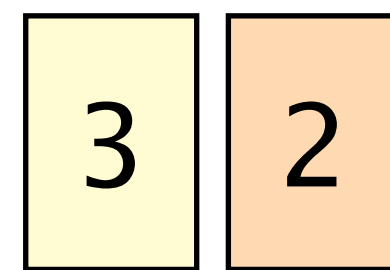


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$

Input sequence

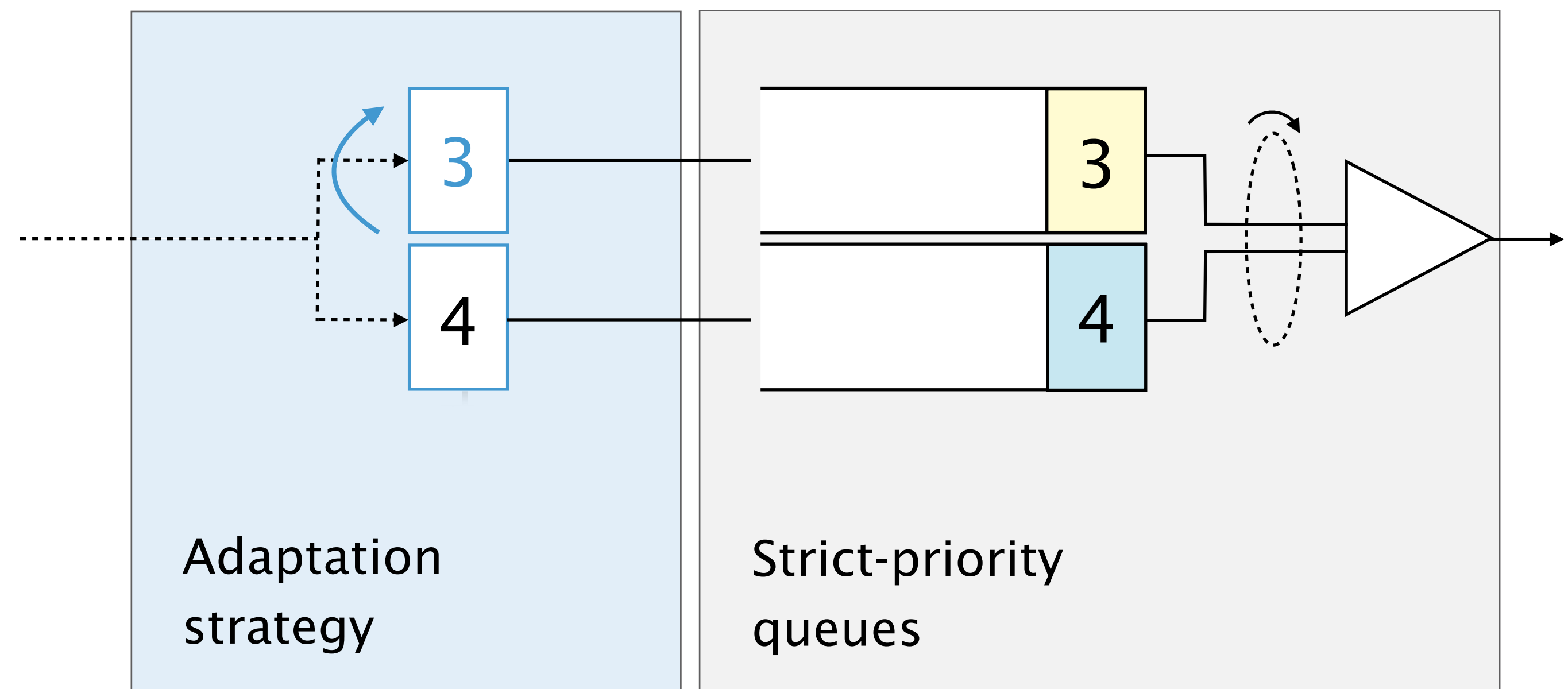
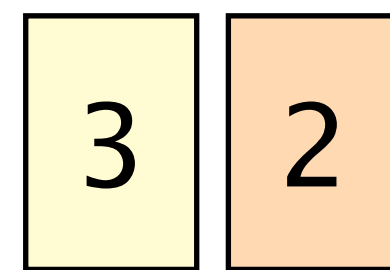


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Push-up adaptation

Set bound to packet rank after enqueue

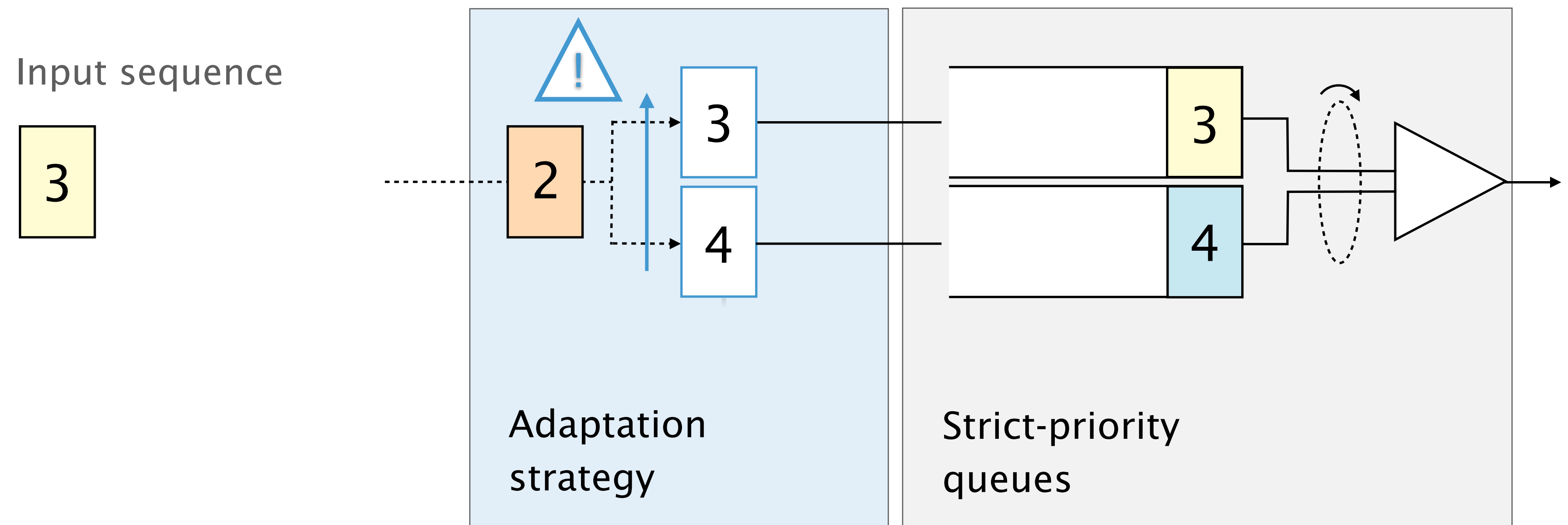
Input sequence



SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

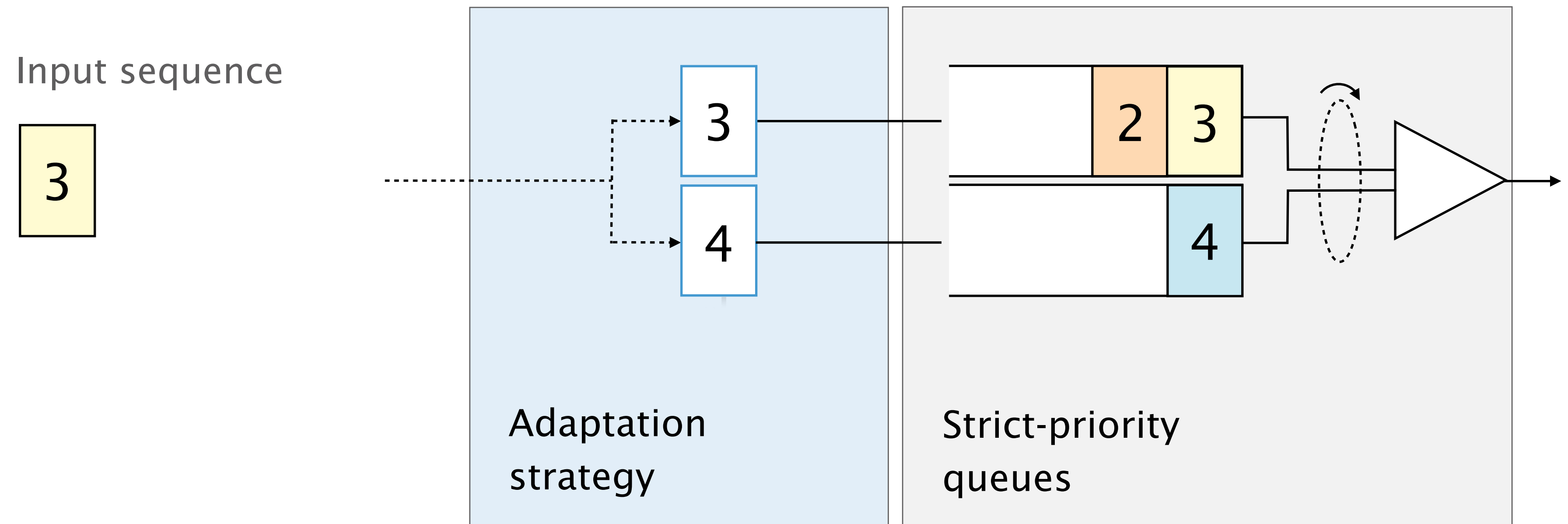
Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$



SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Push-down adaptation

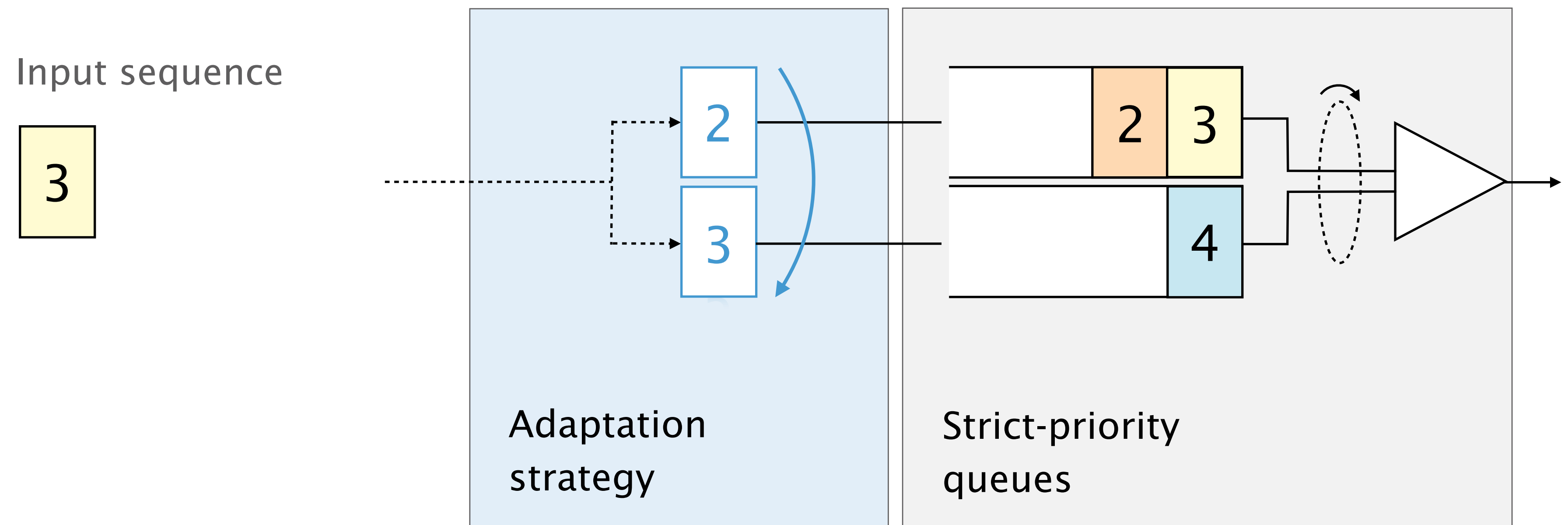
Decrease all bounds after inversion, by inversion cost



SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Push-down adaptation

Decrease all bounds after inversion, by inversion cost

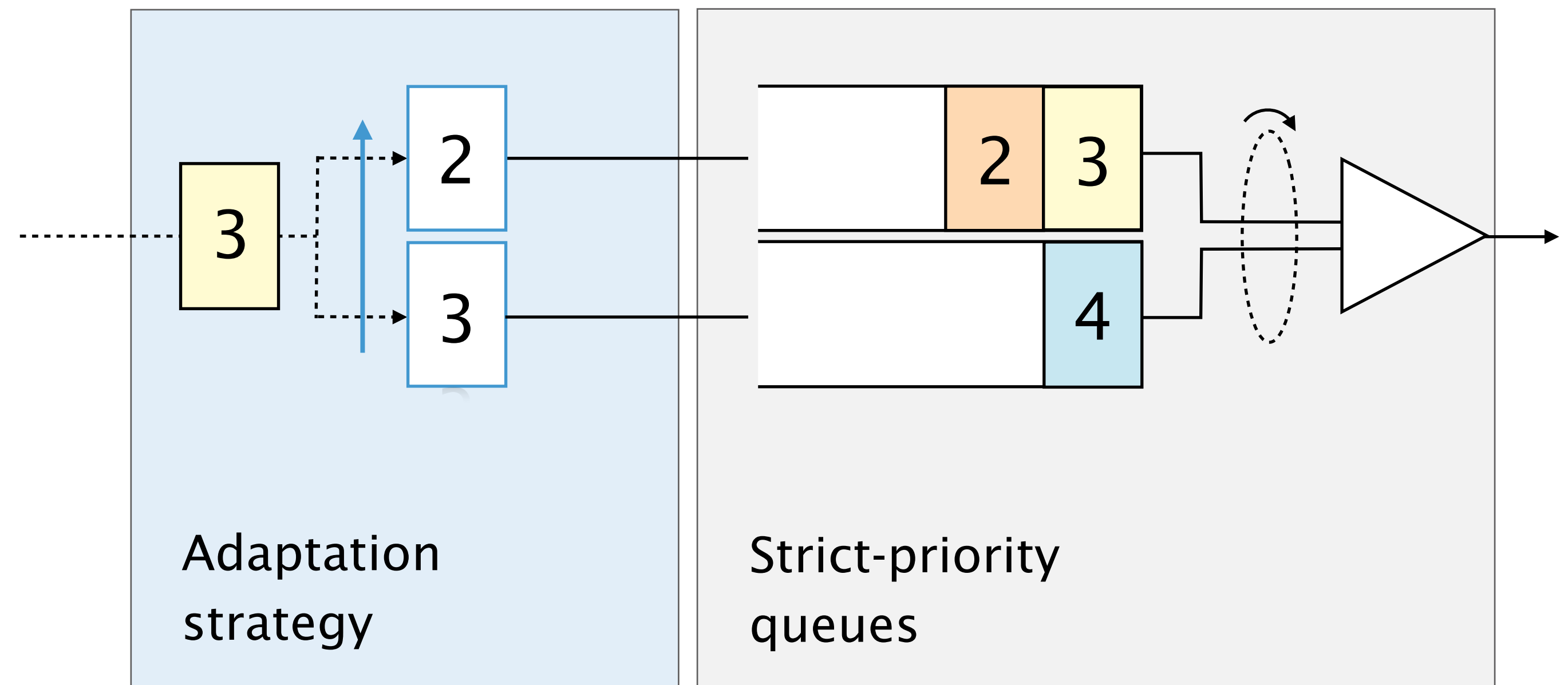


SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$

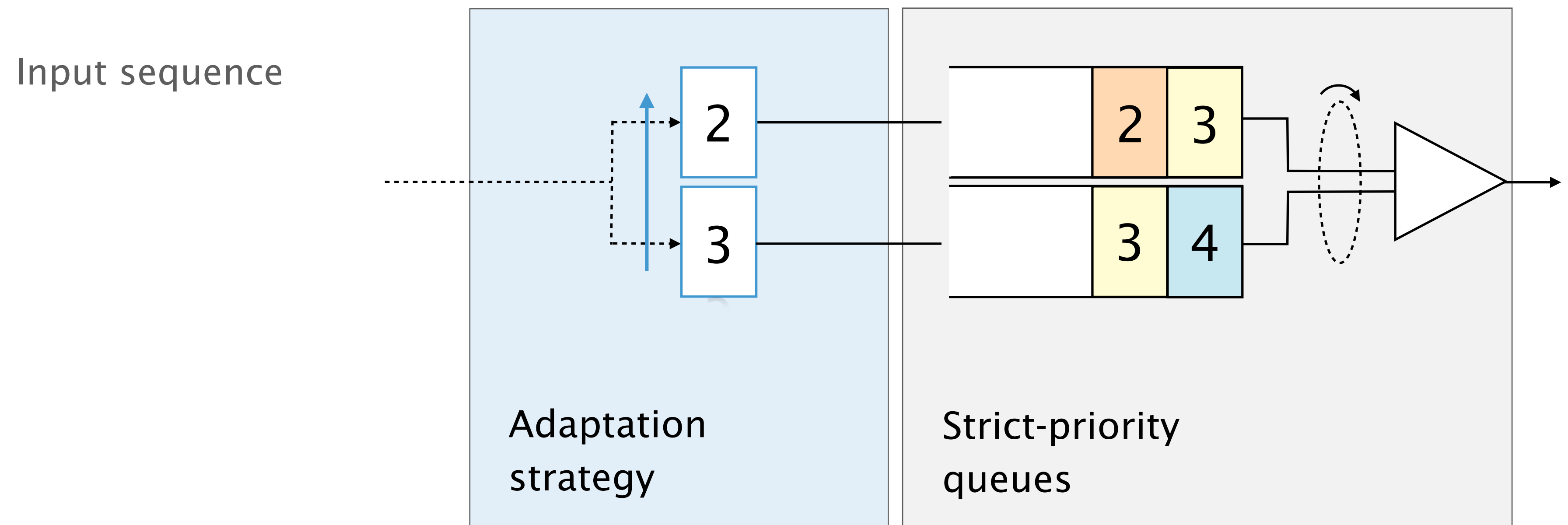
Input sequence



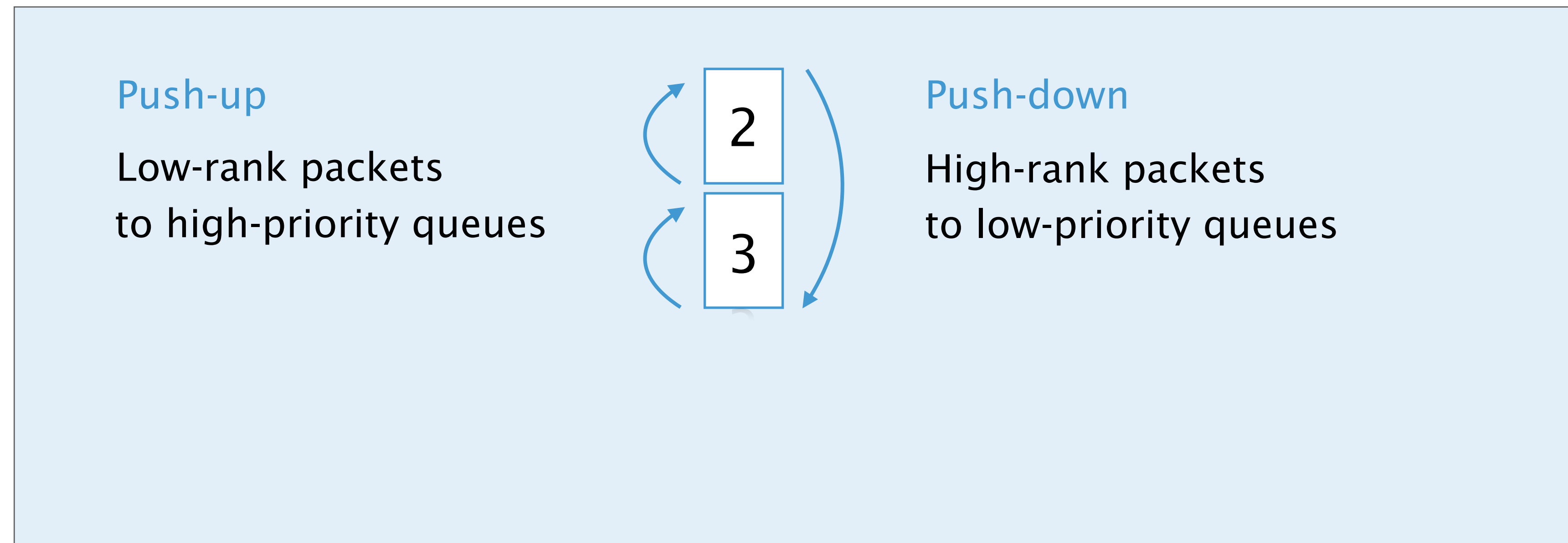
SP-PIFO adapts the mapping of packet ranks to strict-priority queues

Mapping

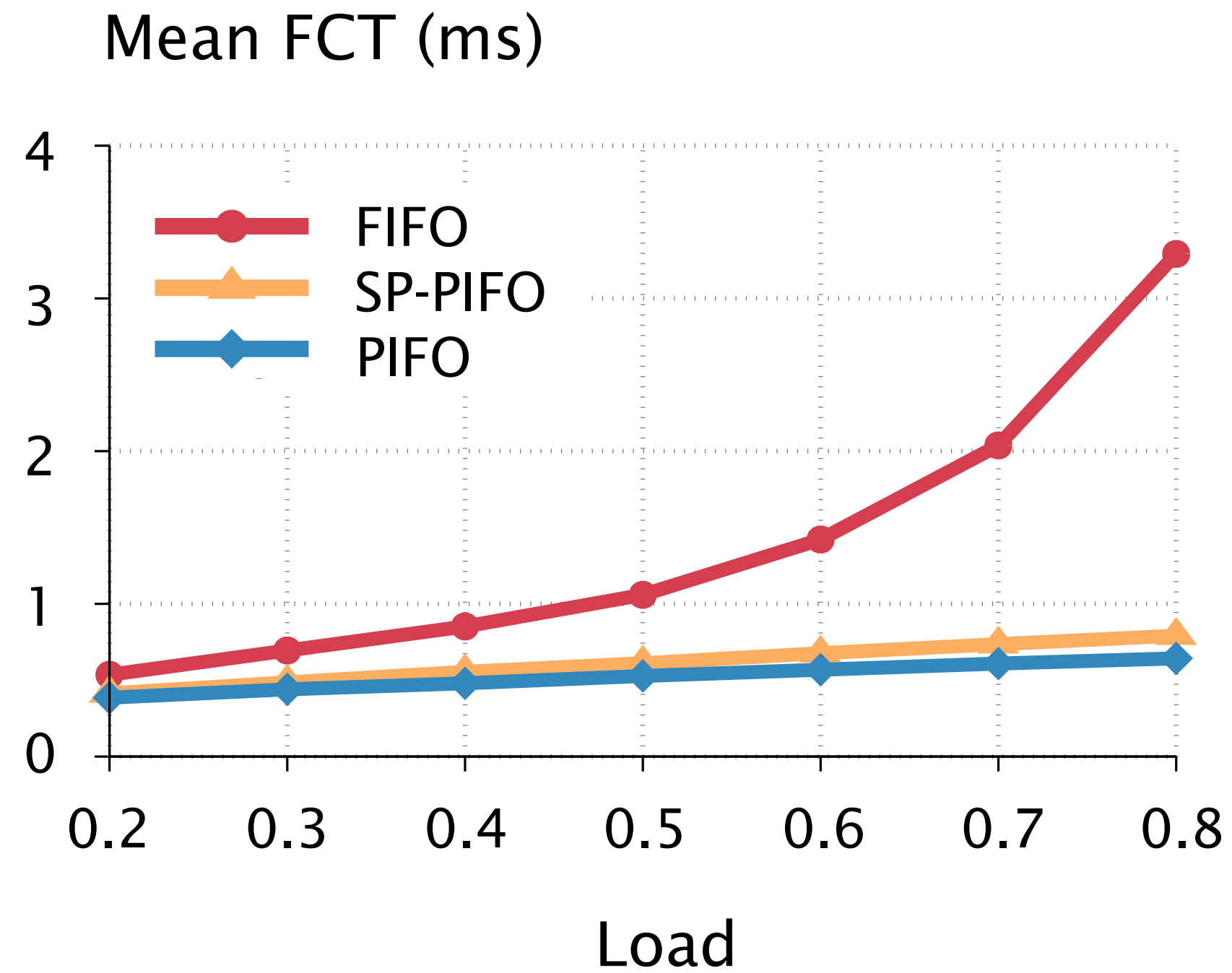
Scan bottom-up, enqueue if $\text{rank} \geq \text{bound}$



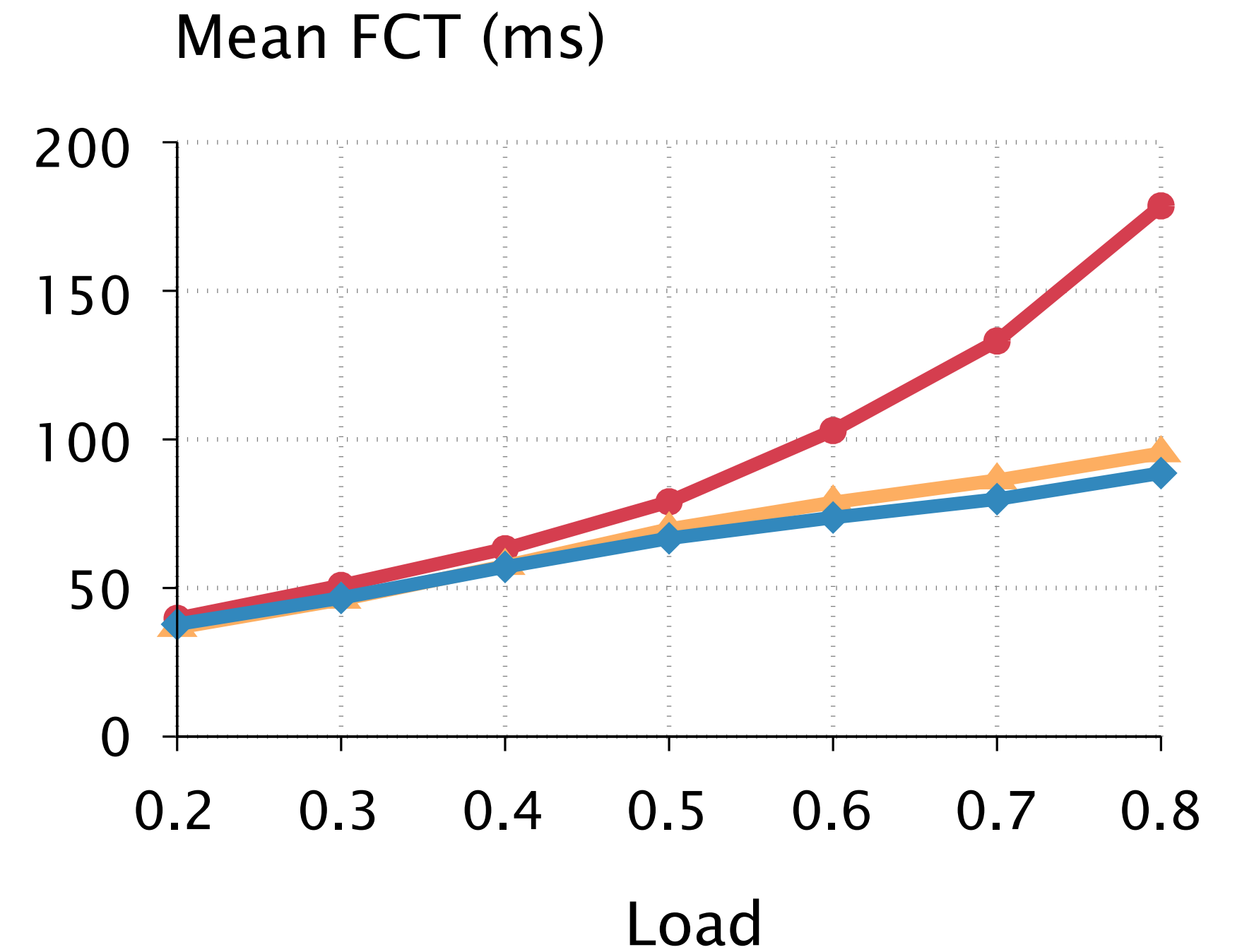
SP-PIFO adapts the mapping of packet ranks to strict-priority queues



SP-PIFO allows us to minimize flow completion times (FCTs)



Small flows <100KB



Big flows ≥ 1 MB

How to enable programmable scheduling on existing devices?

SP-PIFO
[NSDI'20]

Approximating
PIFO's scheduling

PACKS
[NSDI'25]

Incorporating
PIFO's admission

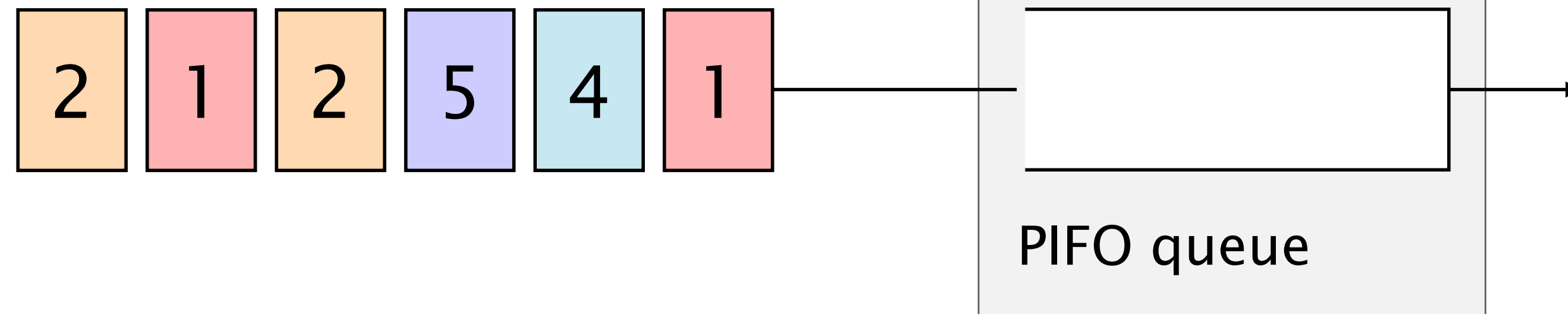
How to use it to improve the Internet's security?

ACC-Turbo
[SIGCOMM'22]

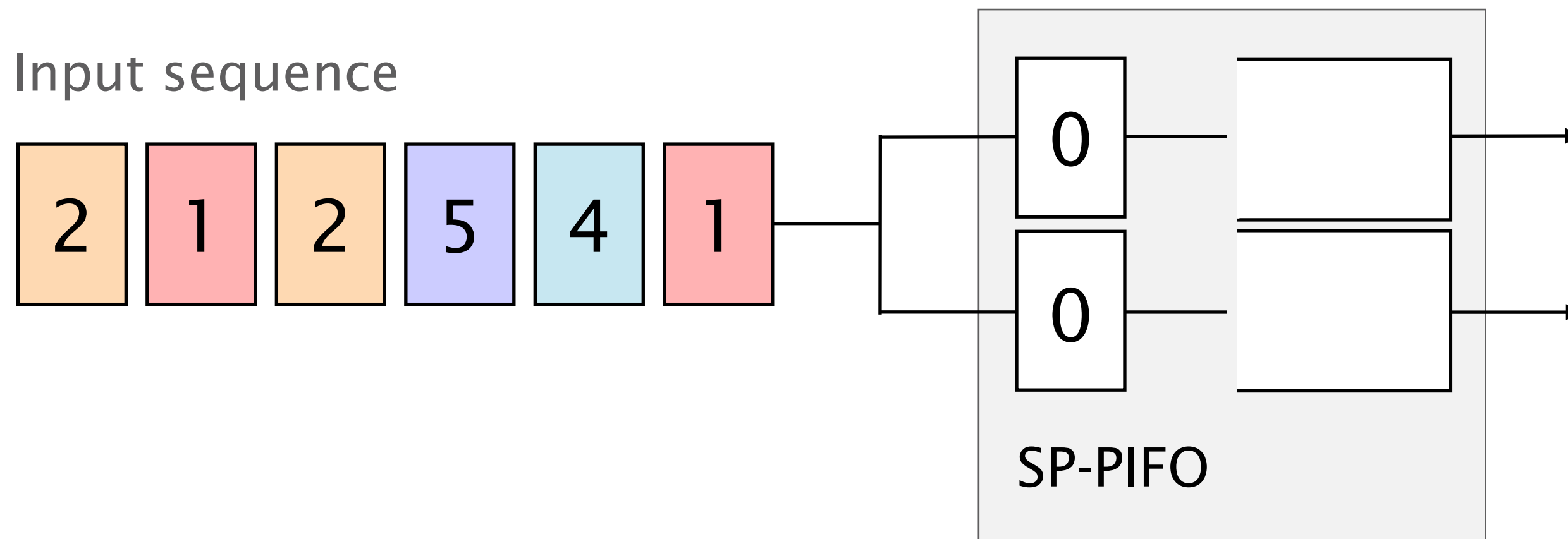
Mitigating
DDoS attacks

PIFO's admission prevents the dropping of important packets

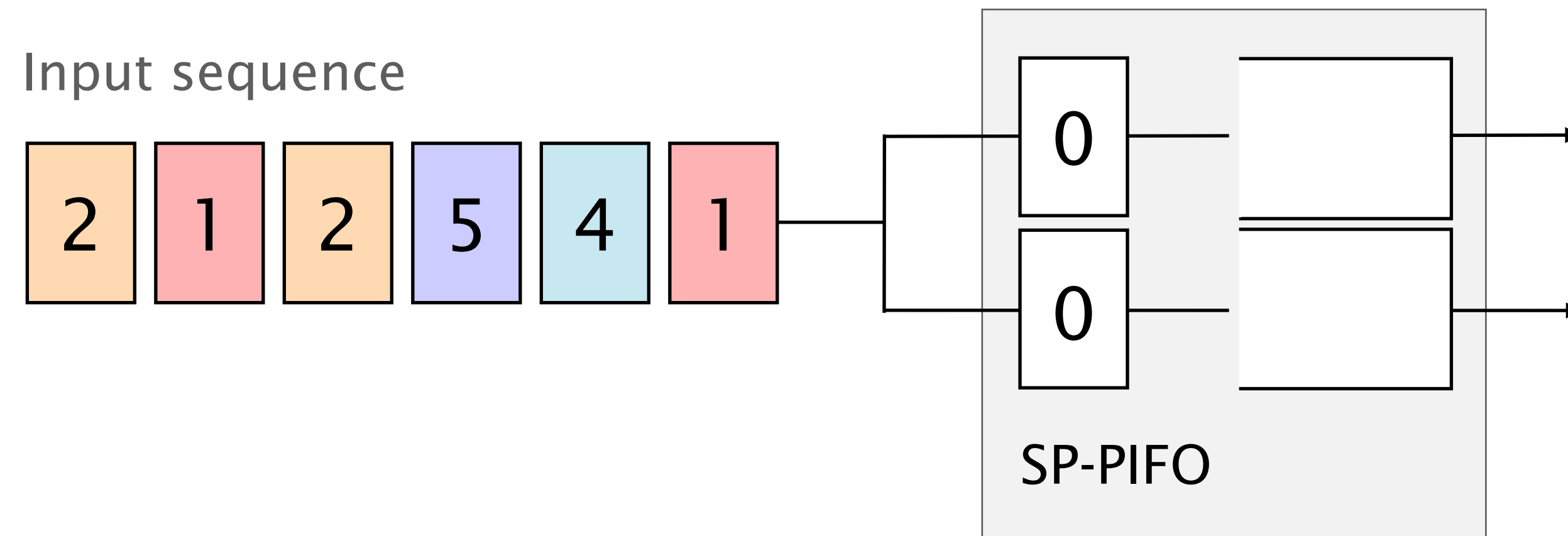
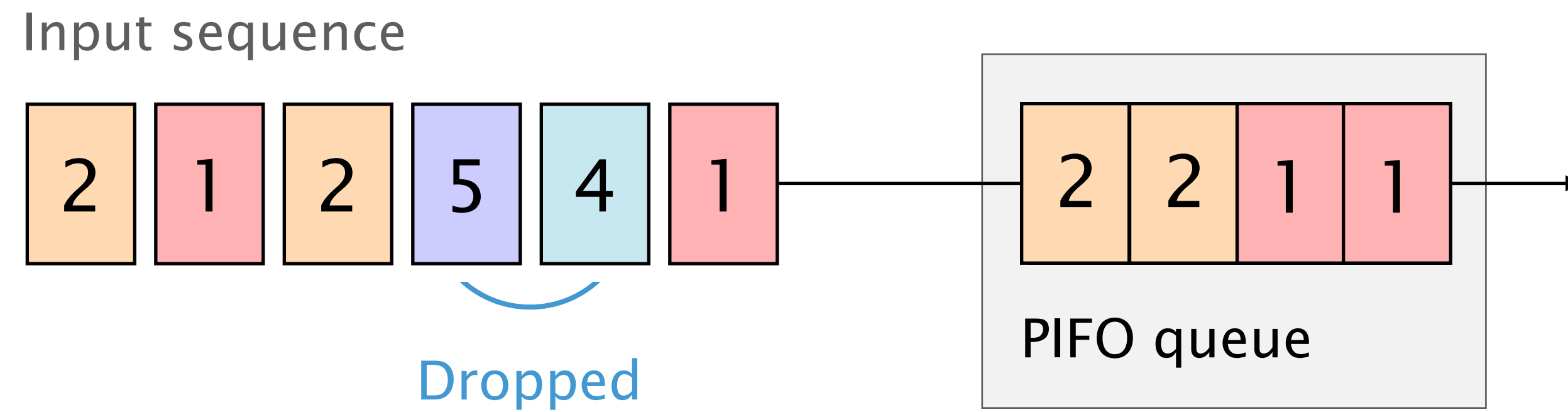
Input sequence



Input sequence

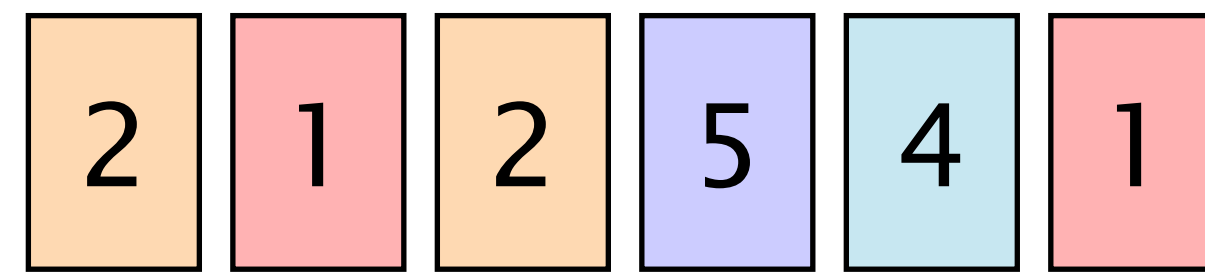


PIFO's admission prevents the dropping of important packets

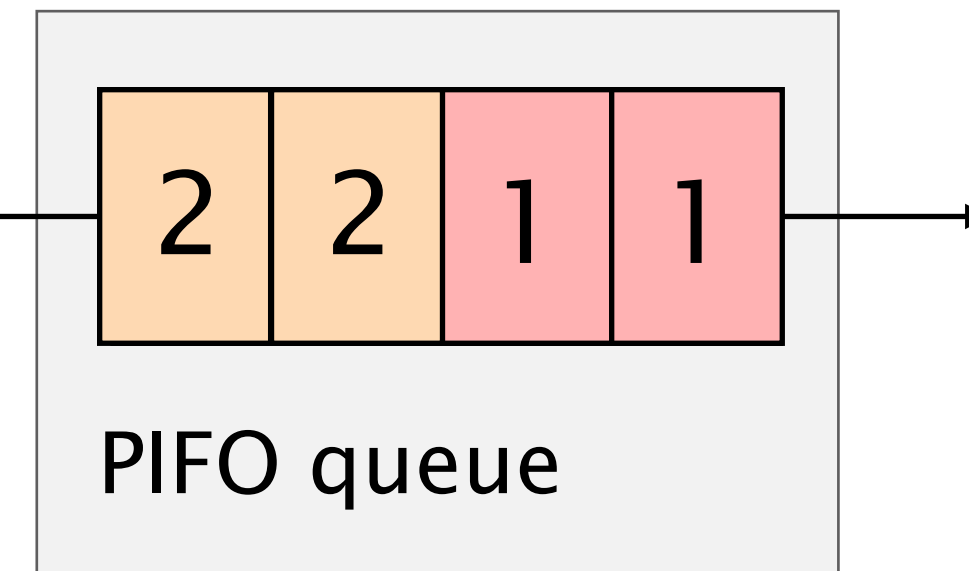


PIFO's admission prevents the dropping of important packets

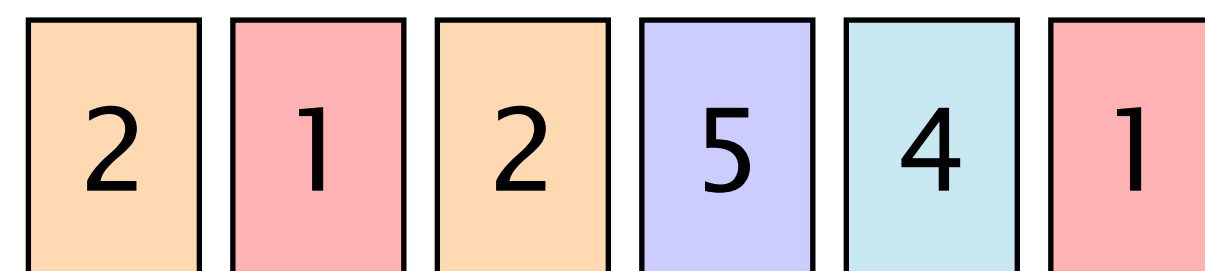
Input sequence



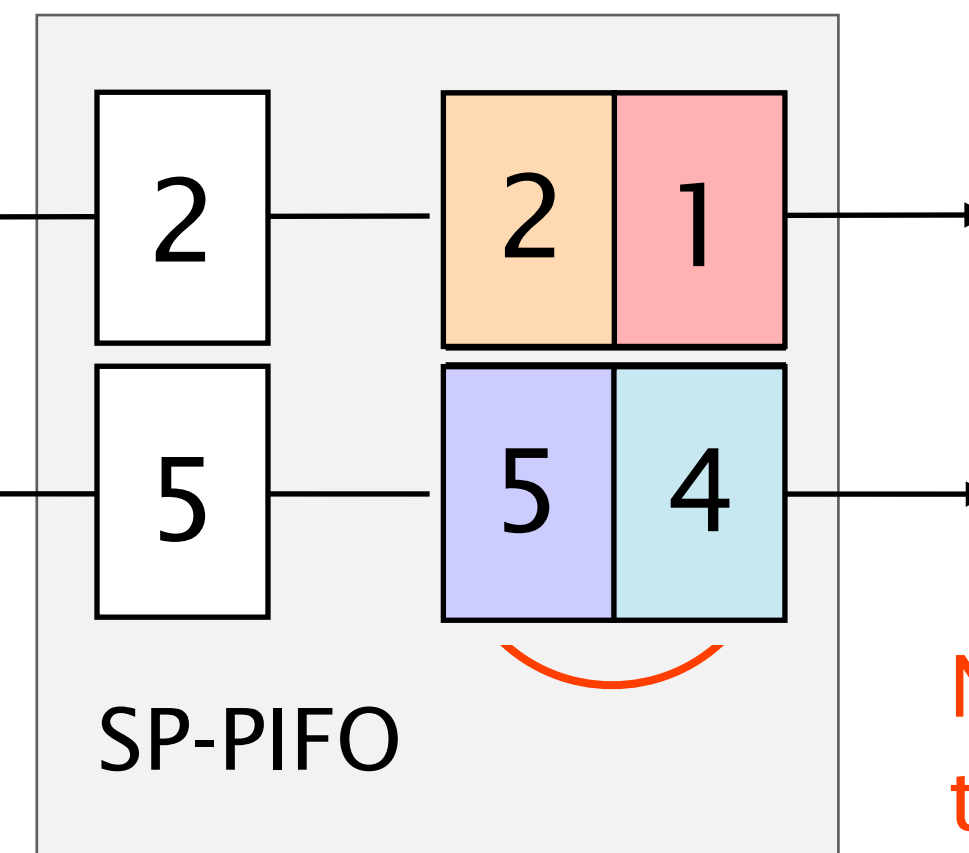
Dropped



Input sequence



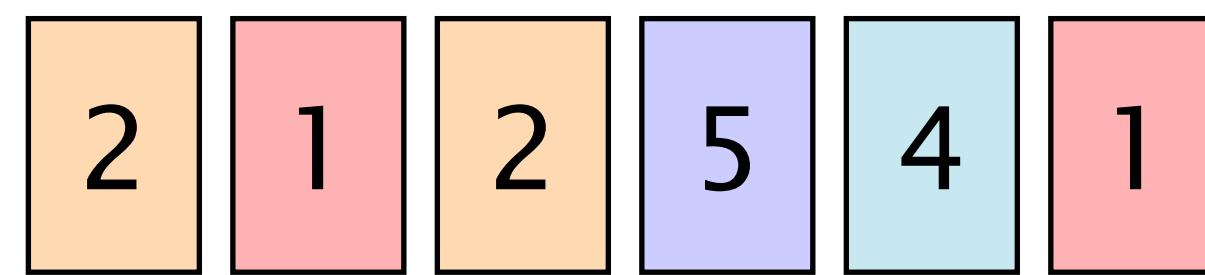
Important packets are dropped



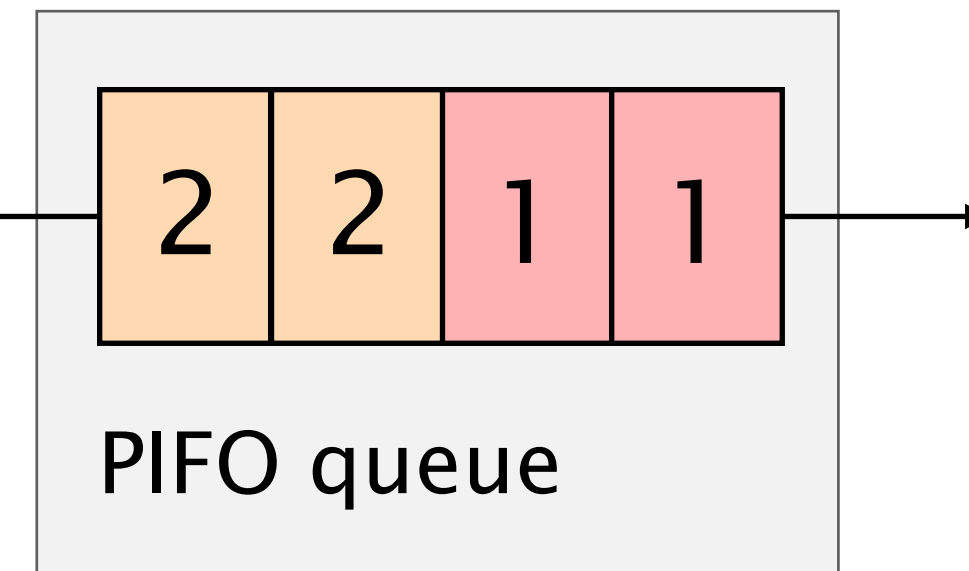
Non-important packets take buffer space

We need to **preemptively** block non-important packets

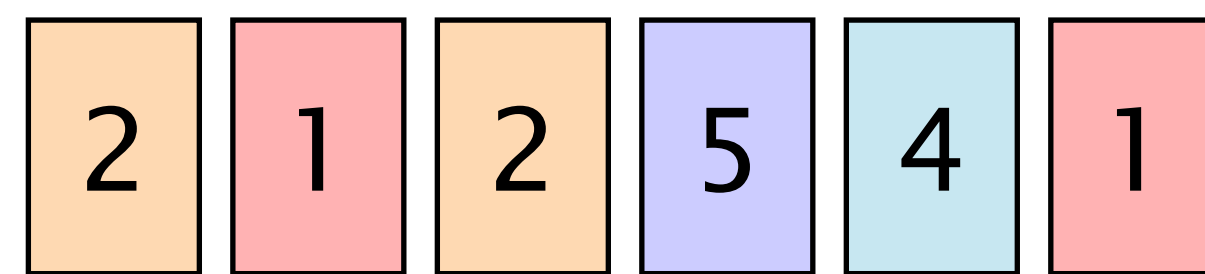
Input sequence



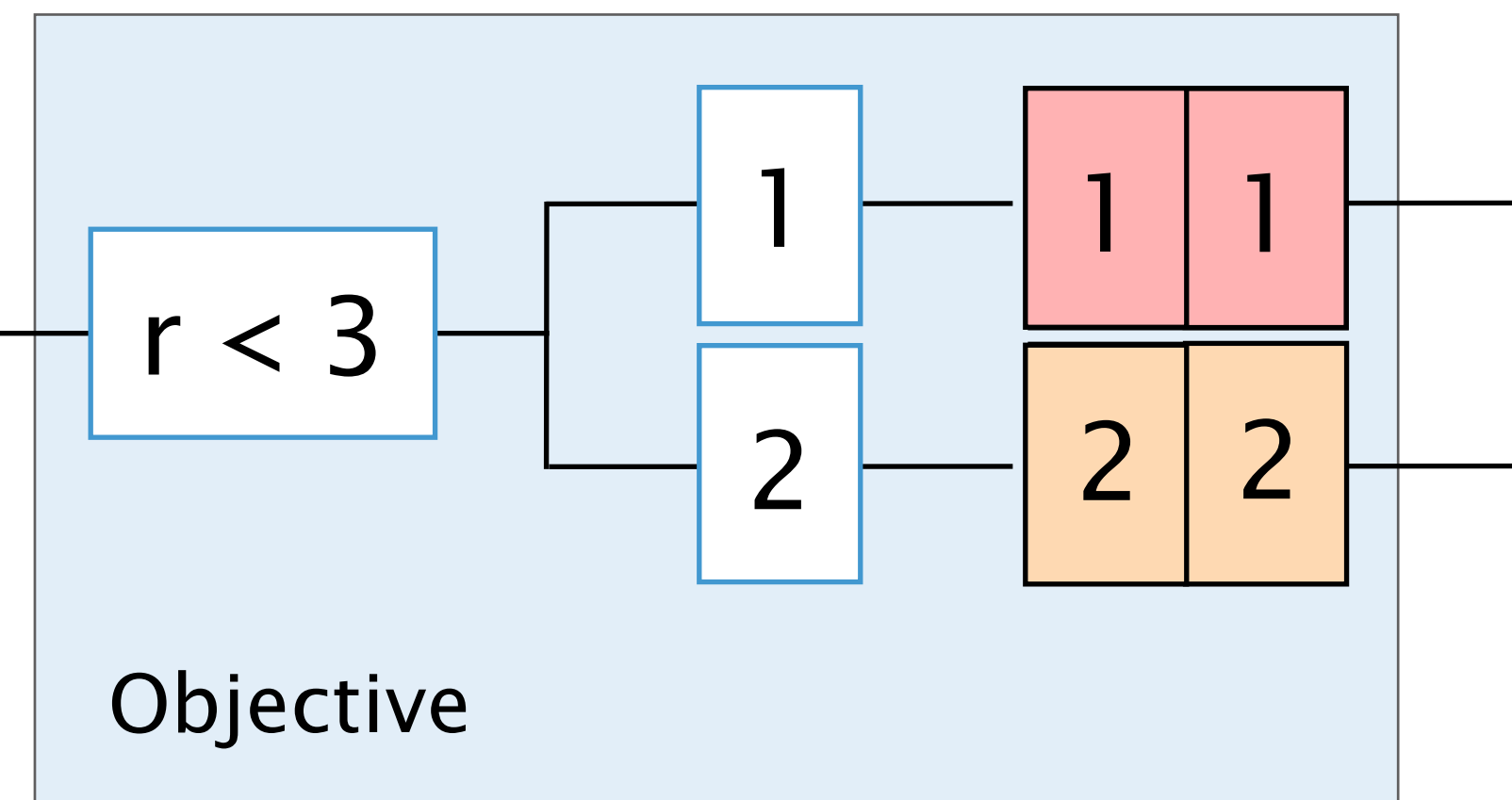
Dropped



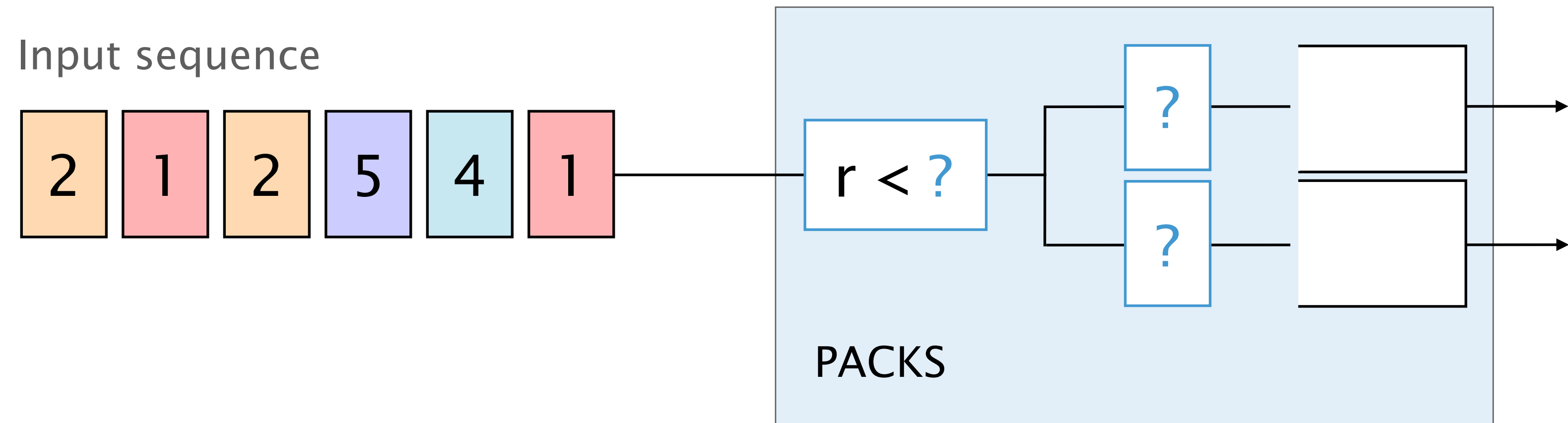
Input sequence



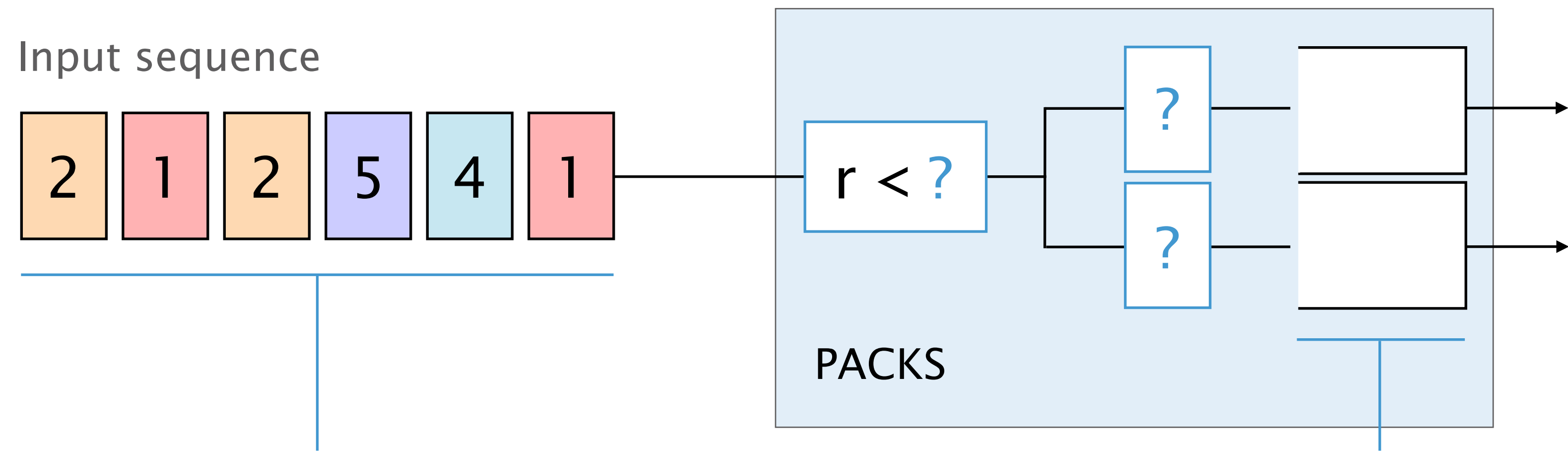
Dropped



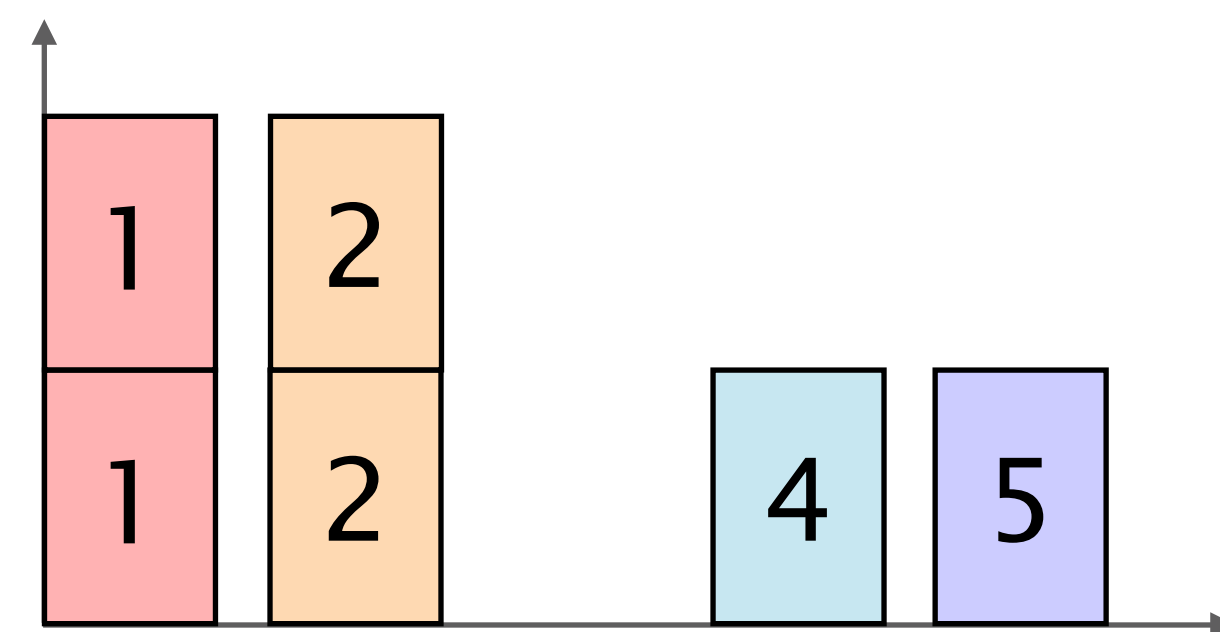
We need to **preemptively** block non-important packets



PACKS monitors the rank distribution and the queue occupancy



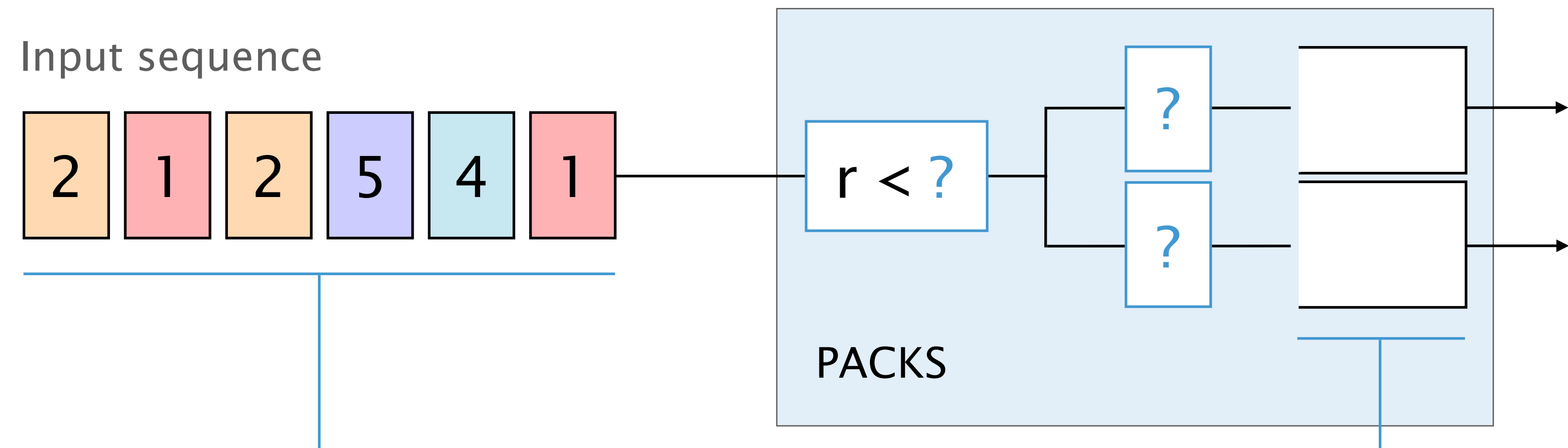
Rank distribution (W)



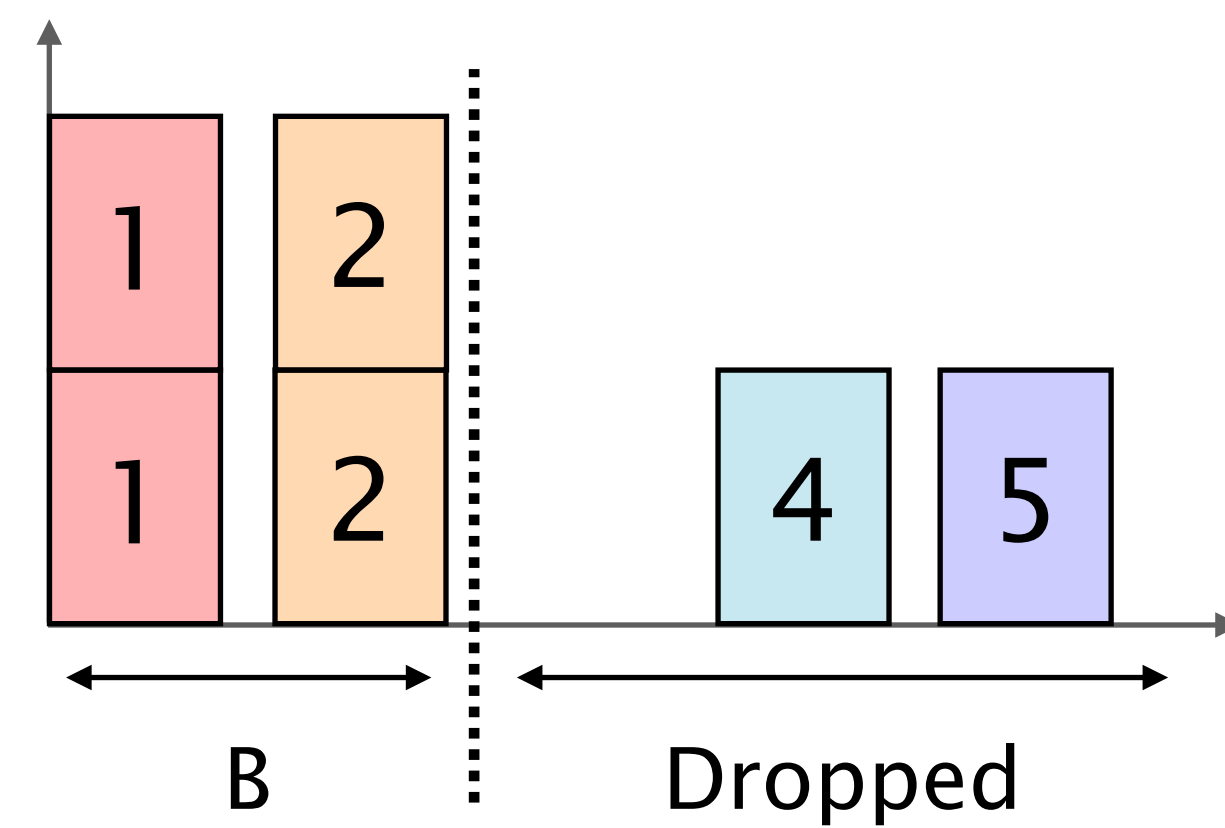
Buffer availability

$B = 4$ packets

PACKS monitors the rank distribution and the queue occupancy



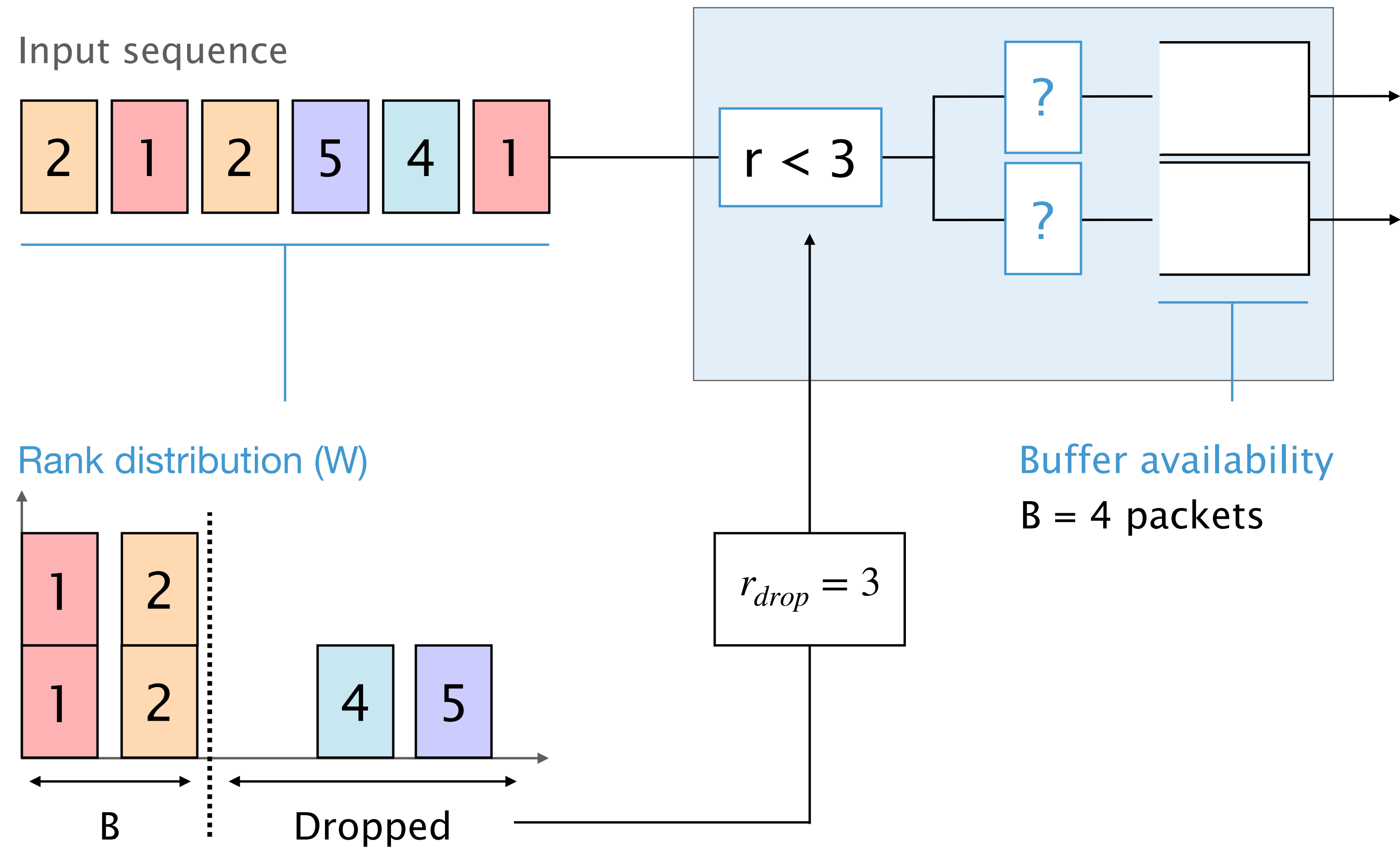
Rank distribution (W)



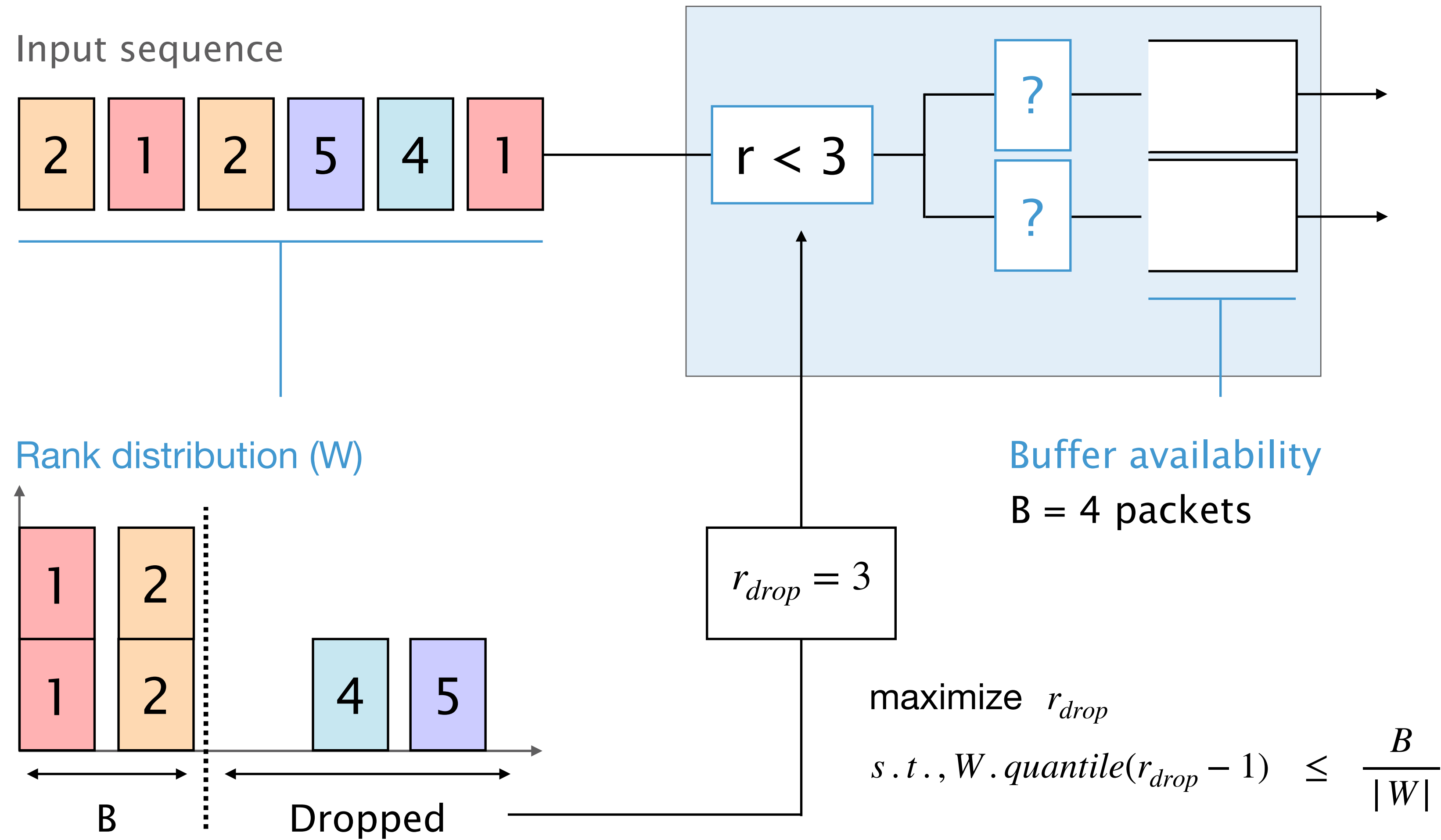
Buffer availability

$B = 4$ packets

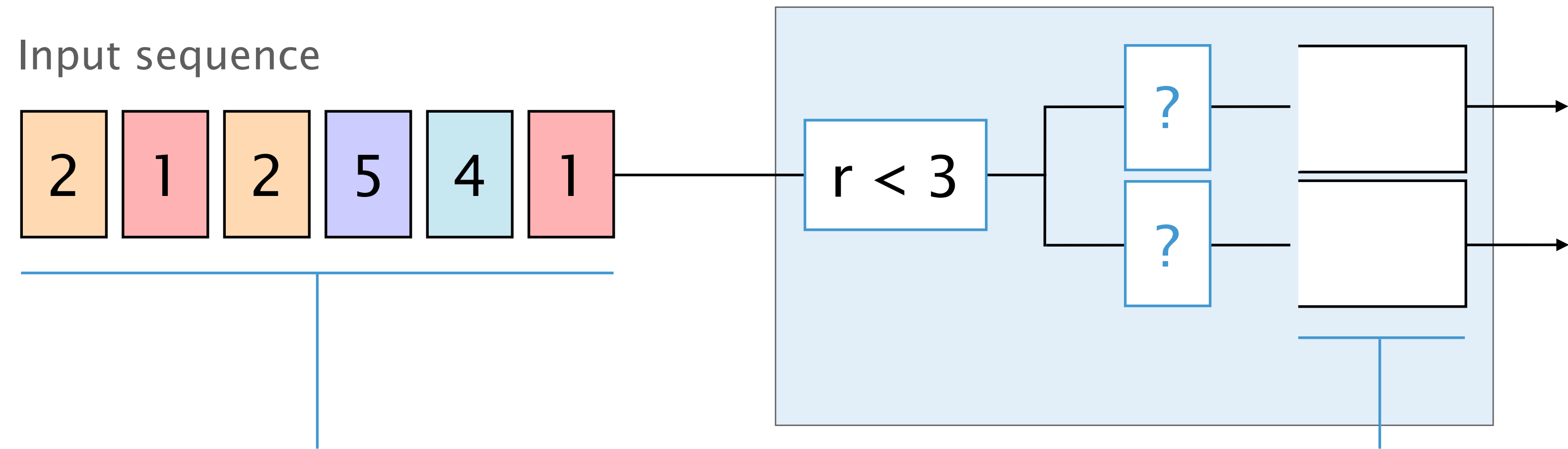
PACKS monitors the rank distribution and the queue occupancy



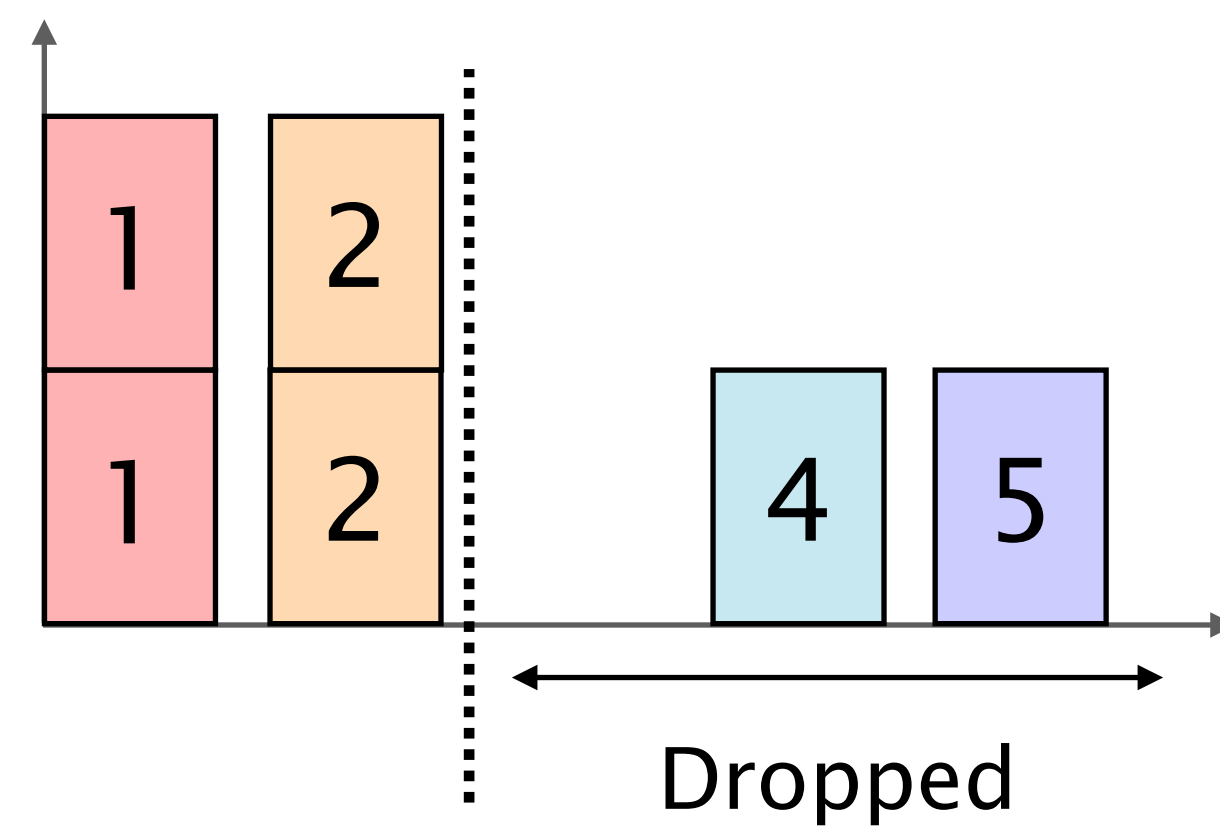
PACKS monitors the rank distribution and the queue occupancy



PACKS monitors the rank distribution and the queue occupancy



Rank distribution (W)

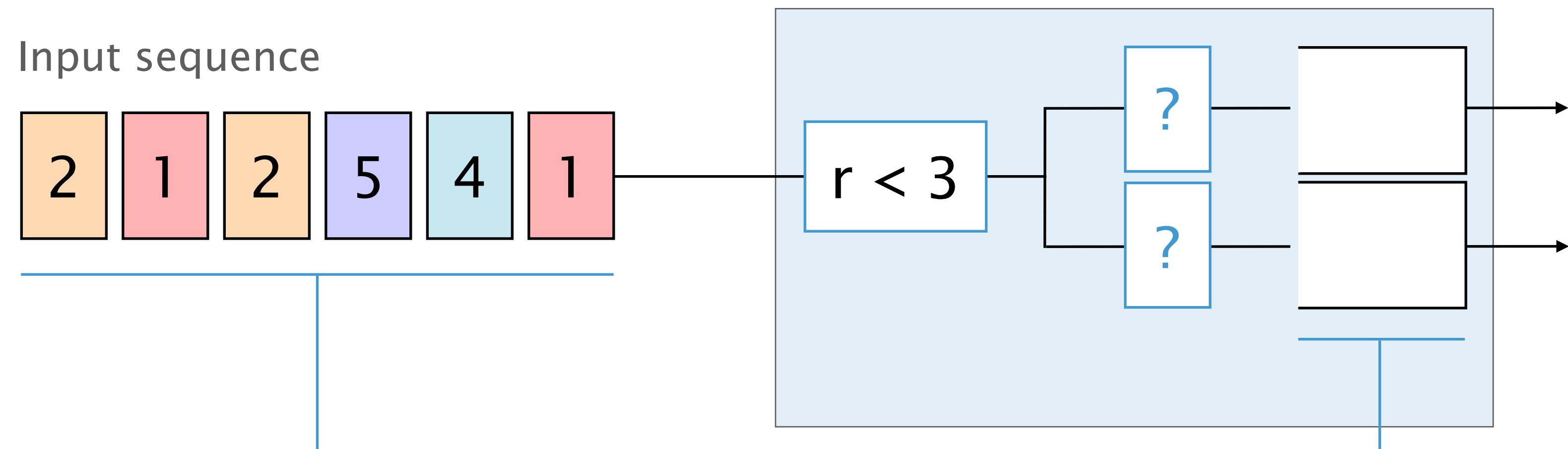


Queues availability

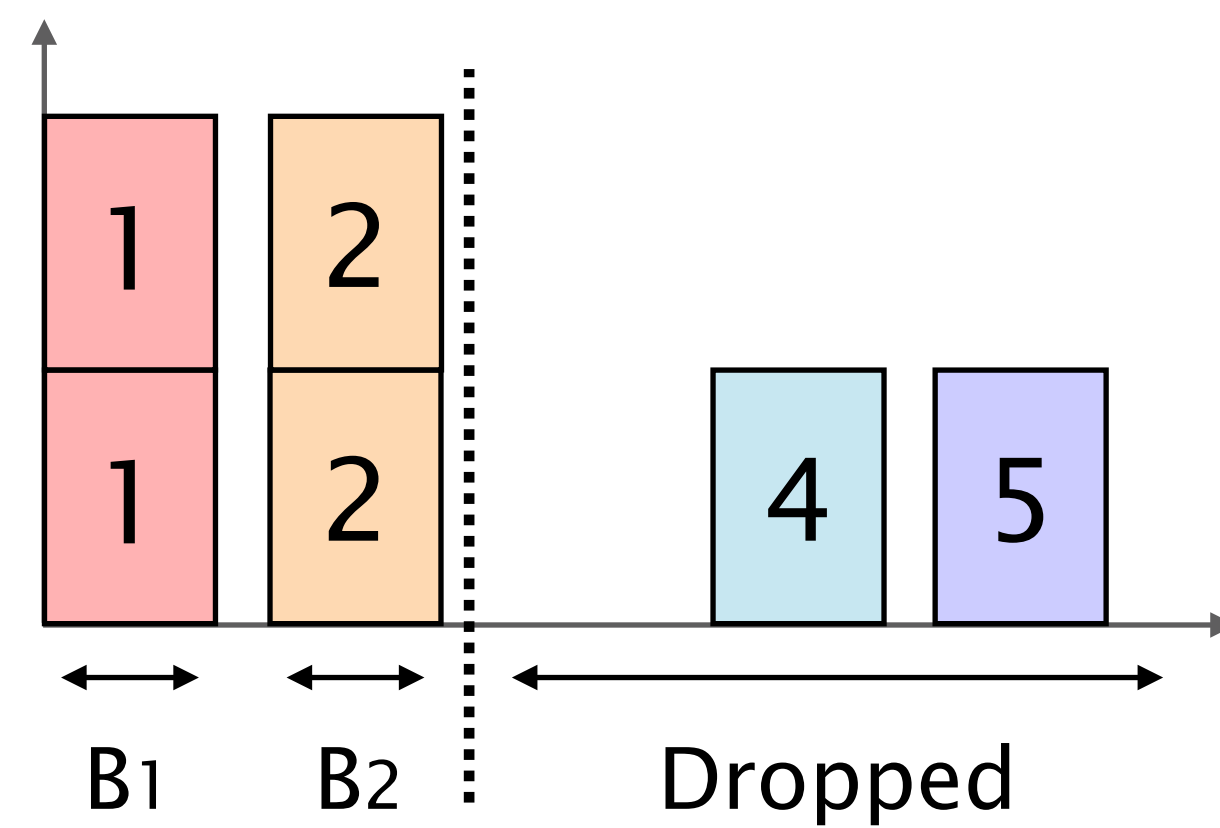
$B_1 = 2$ packets

$B_2 = 2$ packets

PACKS monitors the rank distribution and the queue occupancy



Rank distribution (W)

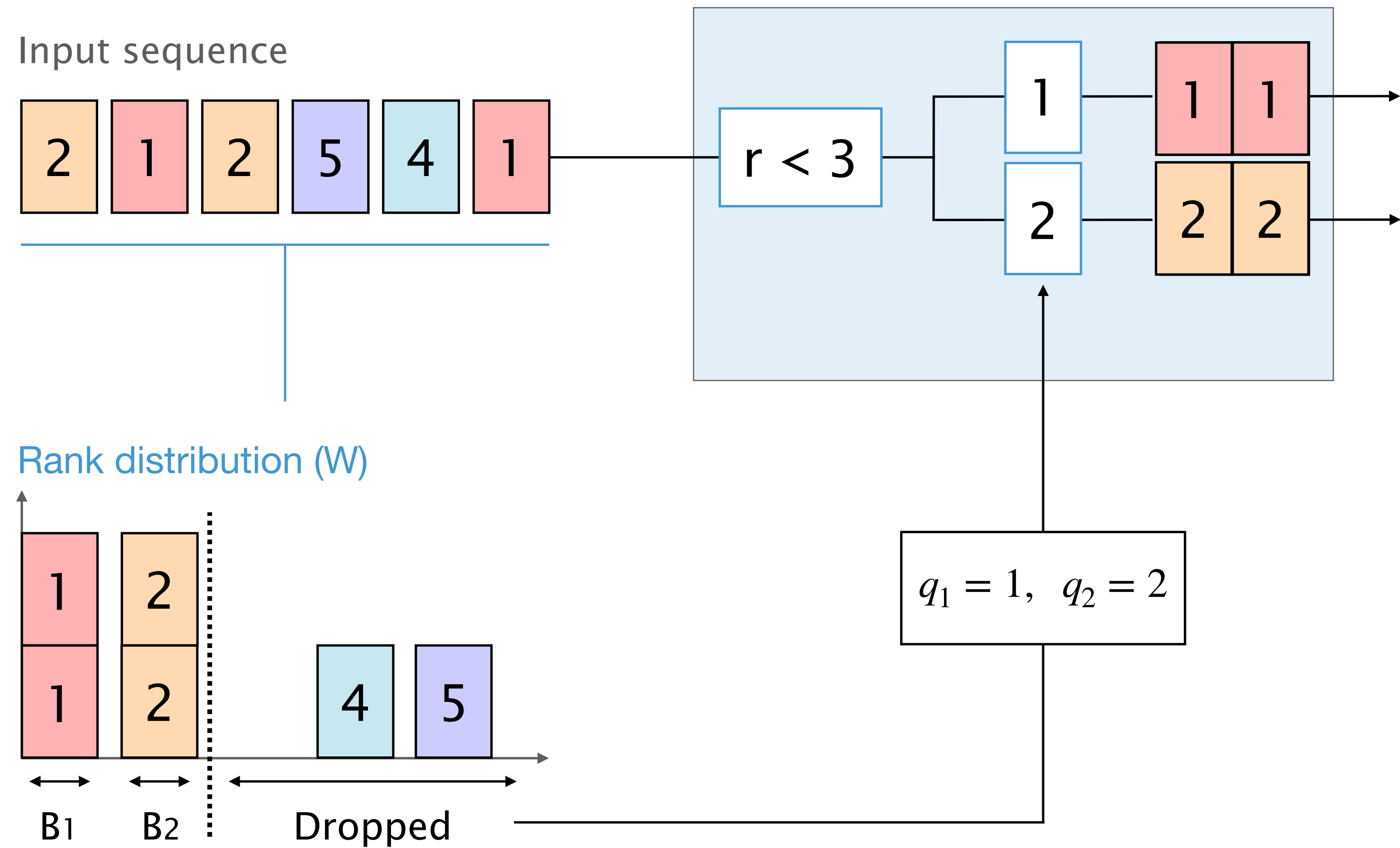


Queues availability

B1 = 2 packets

B2 = 2 packets

PACKS monitors the rank distribution and the queue occupancy



SP-PIFO

Per-packet heuristic

No traffic knowledge

No queue information

PACKS

Window-based

Rank-distribution aware

Queue-occupancy aware

SP-PIFO

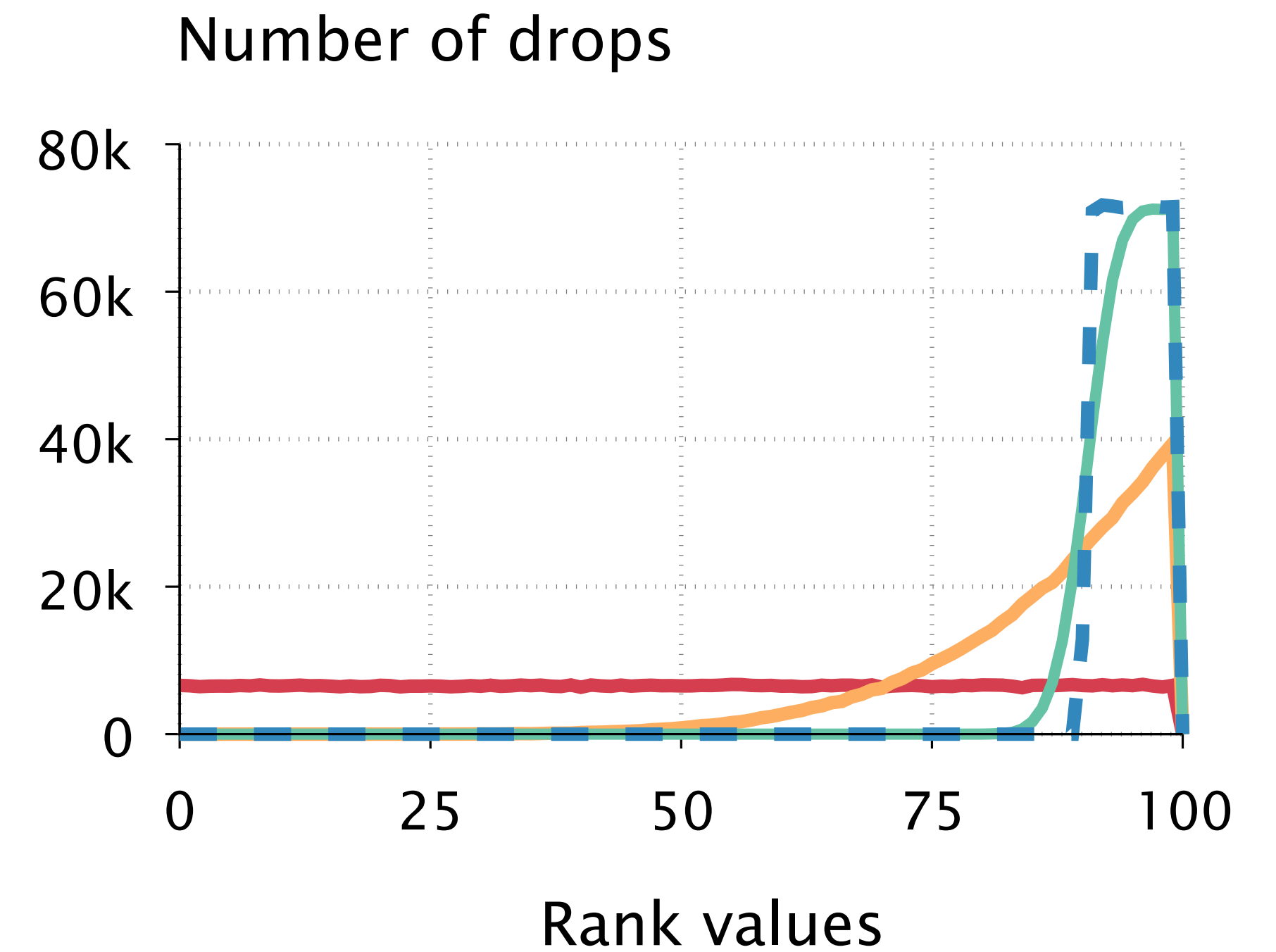
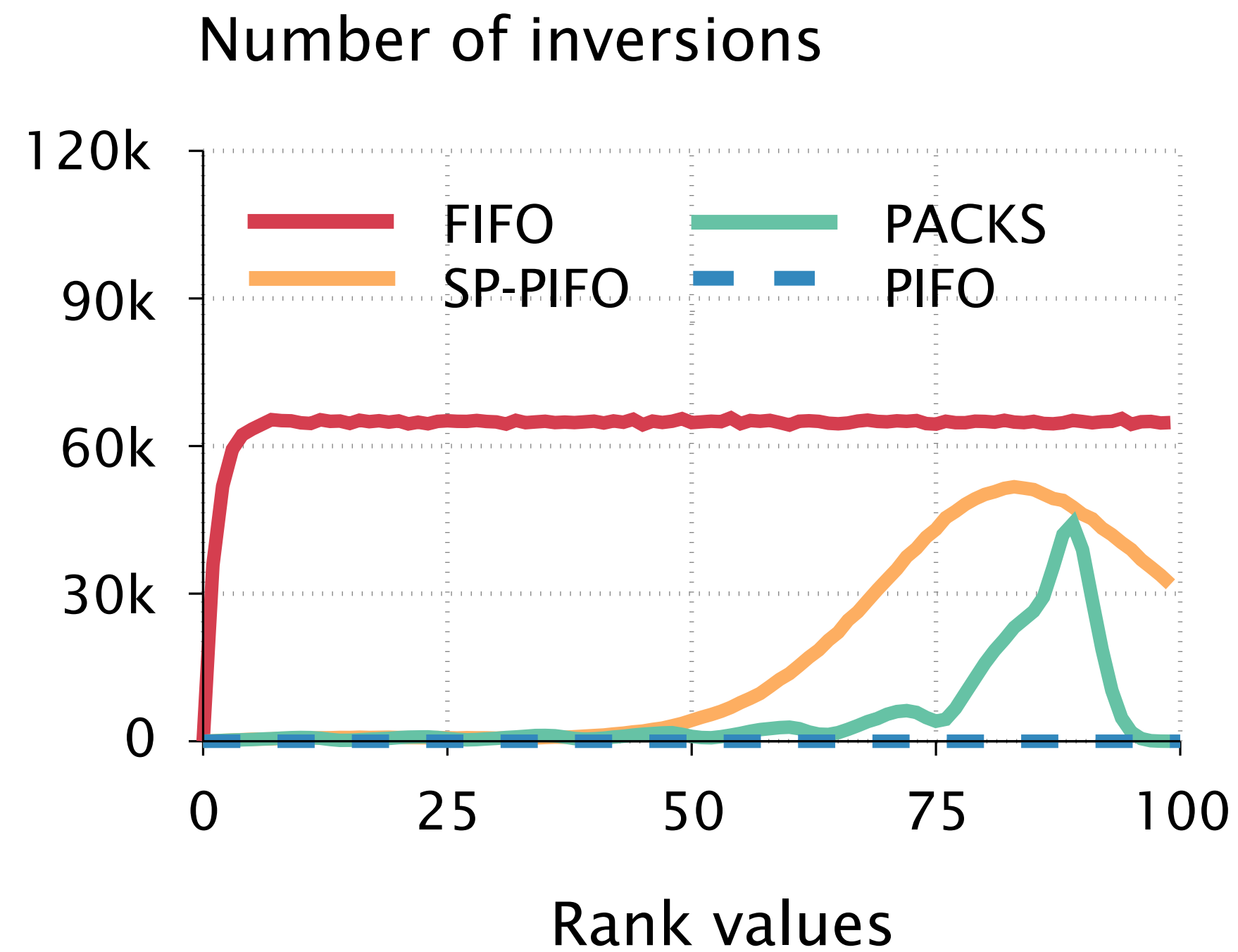
Scheduling ✓

PACKS

Scheduling ✓

Admission ✓

PACKS reduces inversions by up to 7x and drops by up to 60% with respect to SP-PIFO



How to enable programmable scheduling on existing devices?

SP-PIFO

[NSDI'20]

Approximating
PIFO's scheduling

PACKS

[NSDI'25]

Incorporating
PIFO's admission

How to use it to improve the Internet's security?

ACC-Turbo

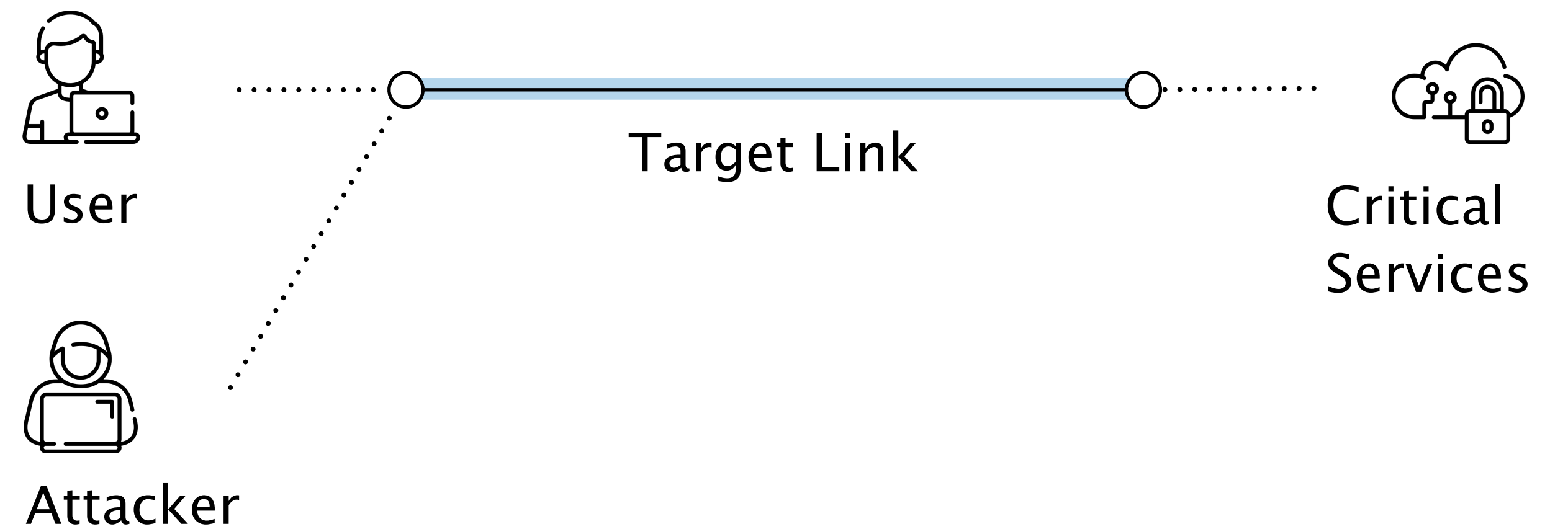
[SIGCOMM'22]

Mitigating
DDoS attacks

Pulse-wave DDoS attacks are
a new type of network-layer DDoS attack

Pulse-wave DDoS attacks are a new type of **network-layer** DDoS attack

Target
a critical link

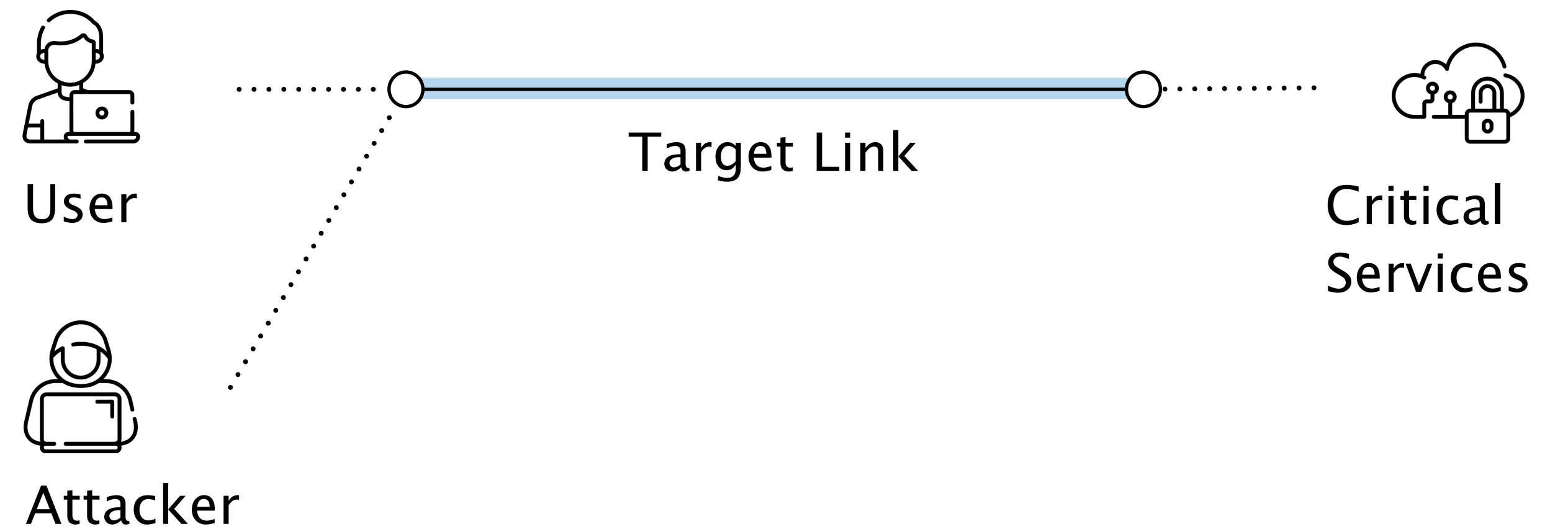


Pulse-wave DDoS attacks are a new type of **network-layer** DDoS attack

Target
a critical link

Volumetric
(Gbps)

Multiple
attack vectors

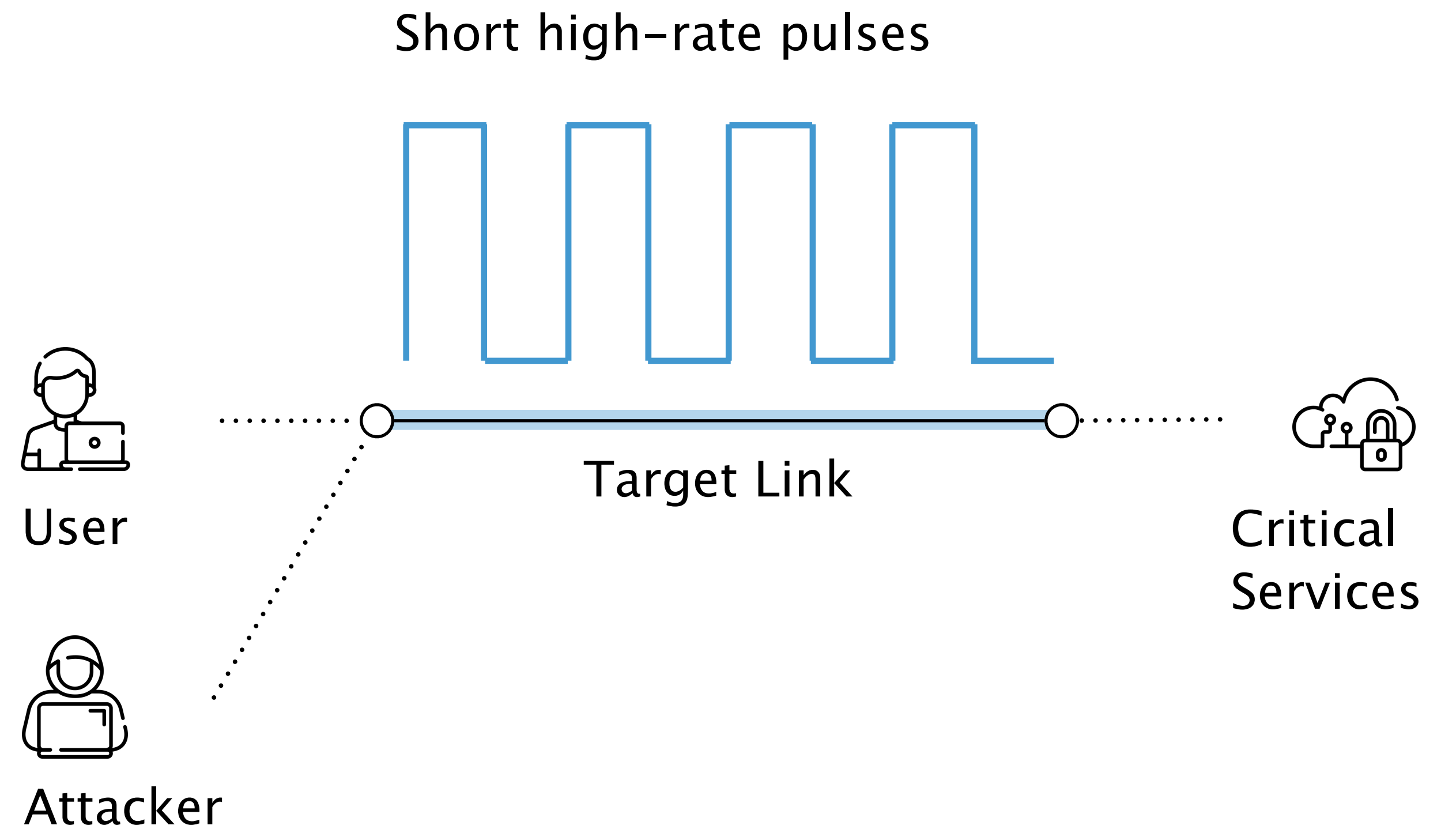


Pulse-wave DDoS attacks are a **new** type of network-layer DDoS attack

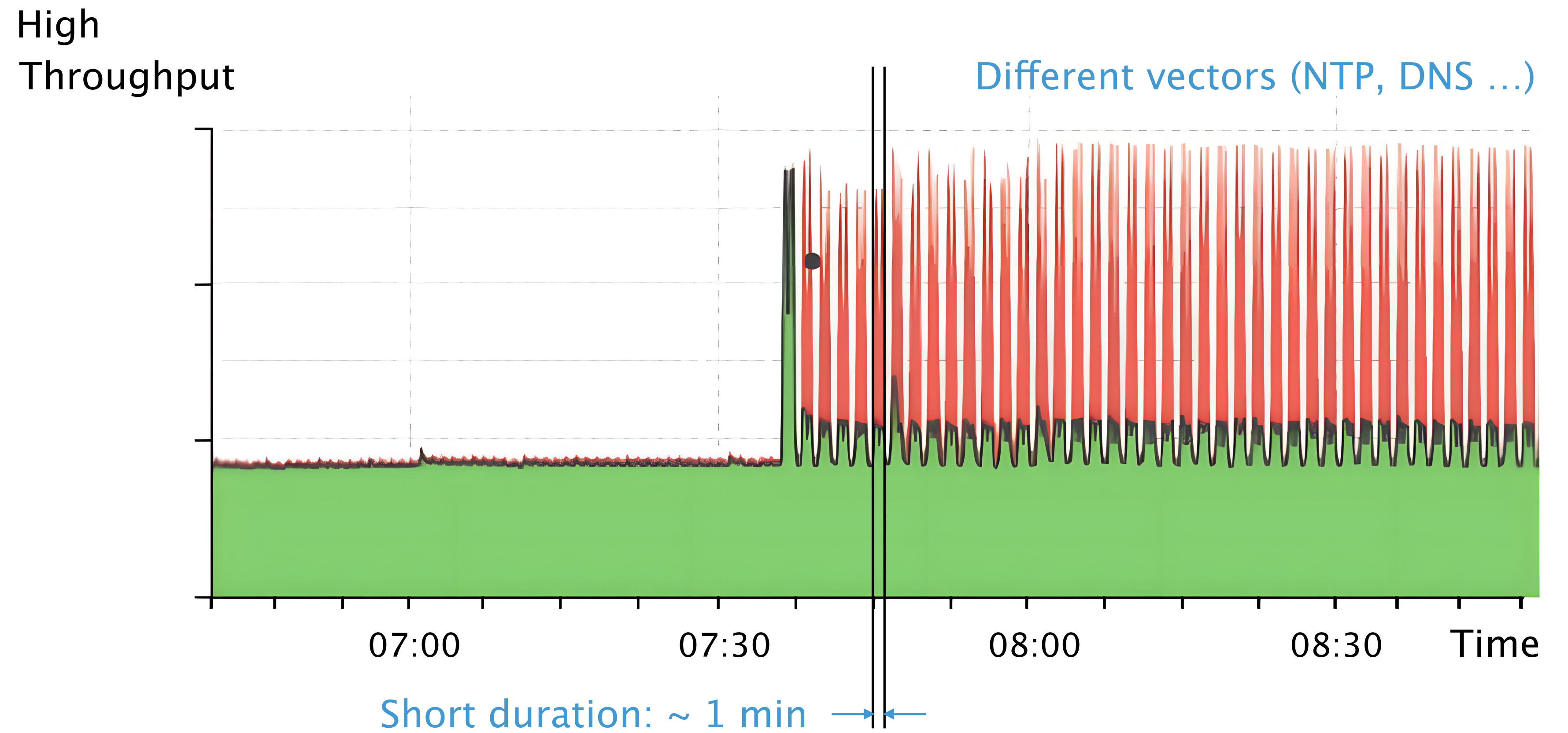
Target
a critical link

Volumetric
(Gbps)

Multiple
attack vectors



Pulse-wave DDoS attacks are an extreme case of congestion

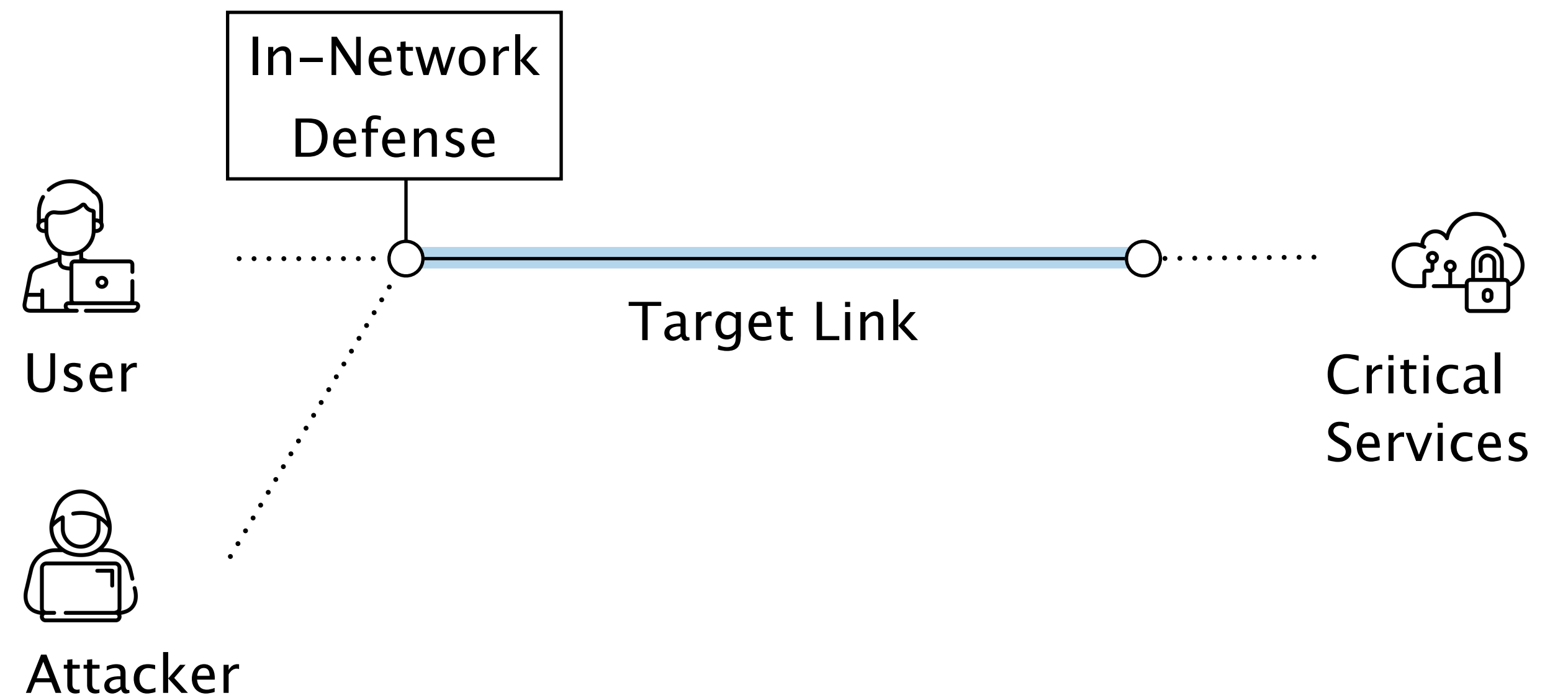


Pulse-wave DDoS attacks exploit
the limitations of existing defenses

Pulse-wave DDoS attacks exploit the **limitations** of existing defenses

Narrow
attack coverage

*Signature-based
Access-control lists*

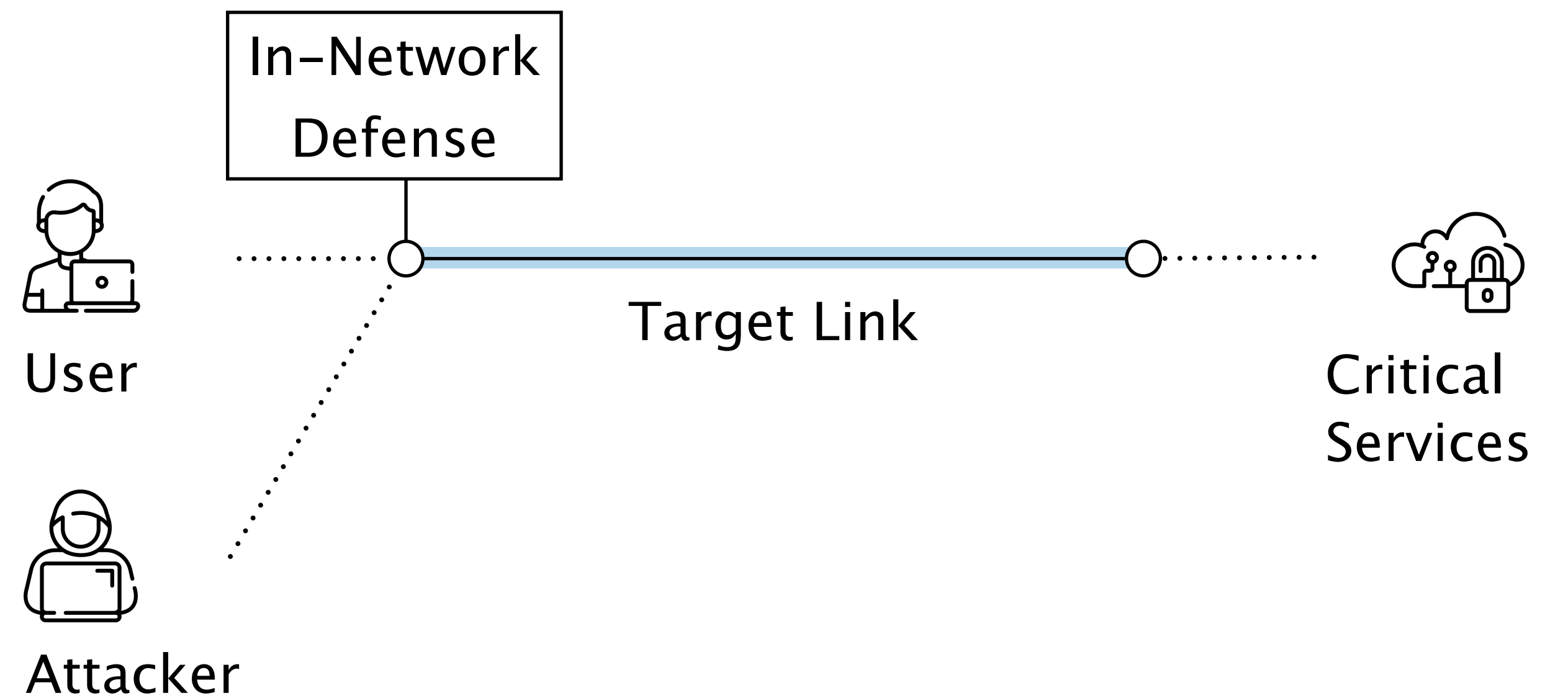


Pulse-wave DDoS attacks exploit the **limitations** of existing defenses

Narrow
attack coverage

Filter-based
Rerouting-based

Drastic
mitigation



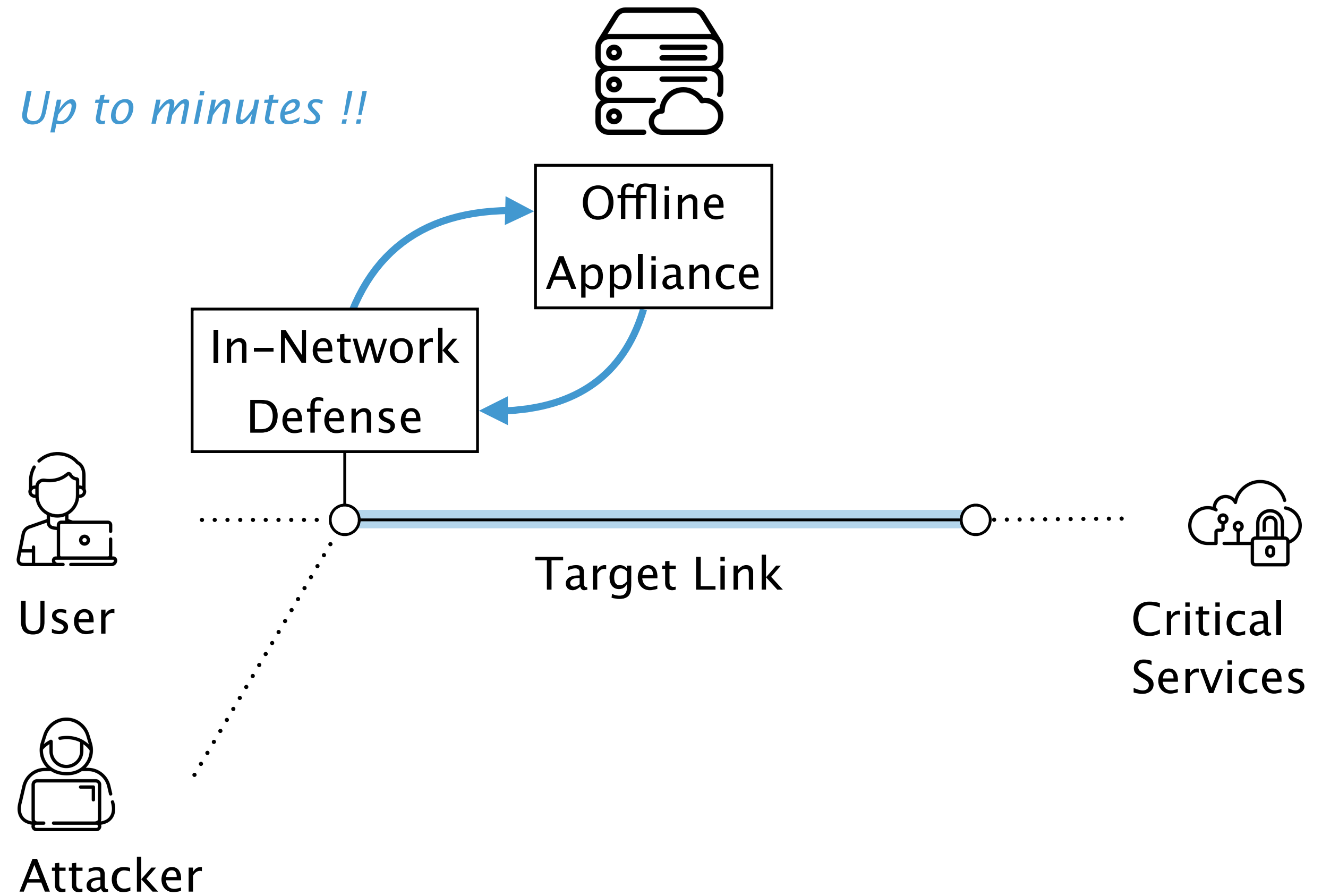
Pulse-wave DDoS attacks exploit the **limitations** of existing defenses

Narrow
attack coverage

Up to minutes !!

Drastic
mitigation

Slow
reaction time



Pulse-wave DDoS attacks exploit the **limitations** of existing defenses

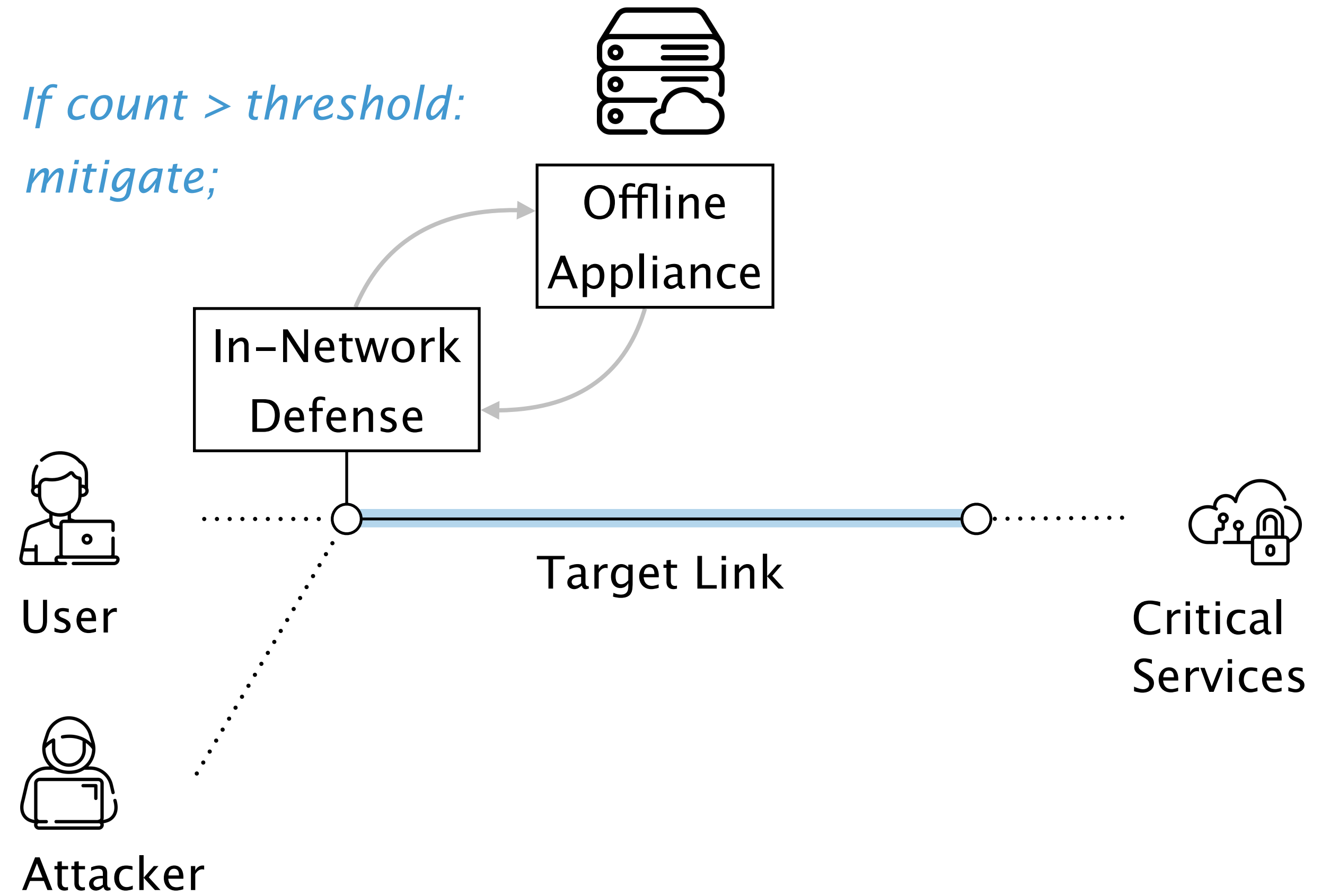
Narrow
attack coverage

Drastic
mitigation

Slow
reaction time

Risk of
misconfiguration

*If count > threshold:
mitigate;*



A pulse-wave DDoS defense needs to be ...

Fast
reaction

Generic
detection

A pulse-wave DDoS defense needs to be ...

Fast
reaction

In-network, at line rate
with limited resources

Generic
detection

Unsupervised techniques
with uncertainty

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Filtering ✘

Rerouting ✘

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Programmable
scheduling

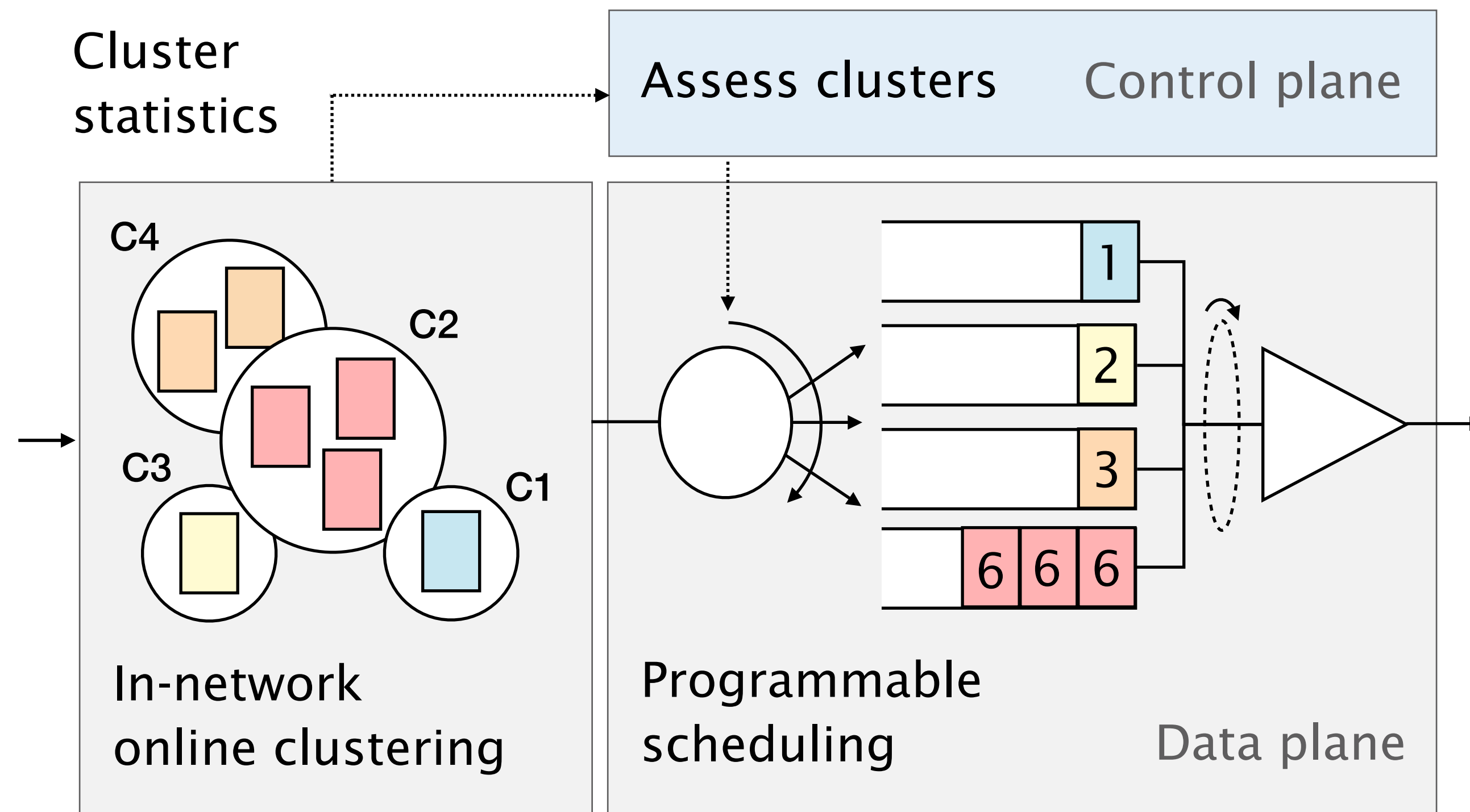
Programmable scheduling is a **safe** mitigation technique

Leverages the whole uncertainty spectrum
with fine-grained scheduling policies

Only drops under congestion
starting by most-malicious packets

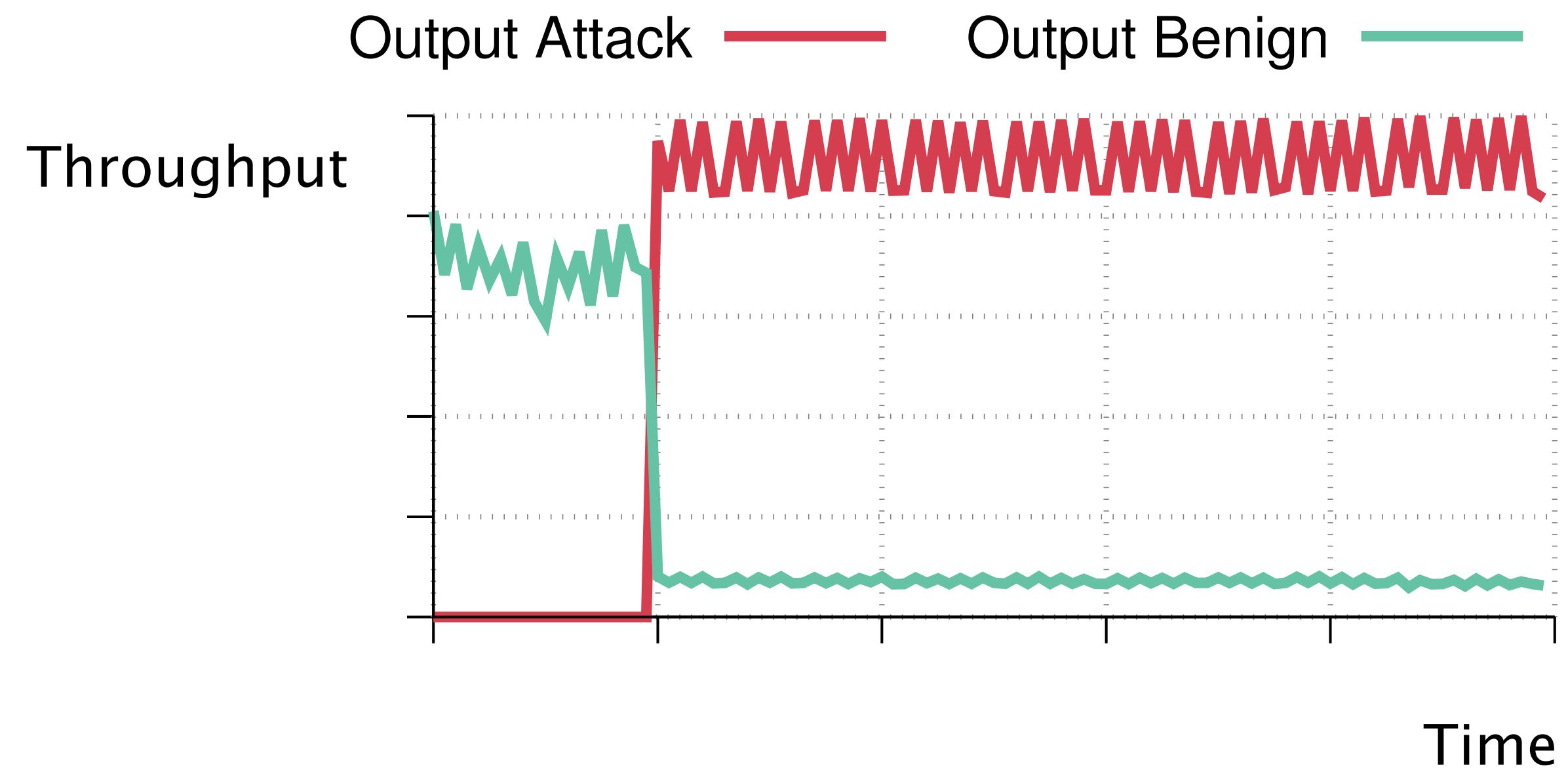
Does not require activation
can be always-on

ACC-Turbo utilizes online clustering and programmable scheduling



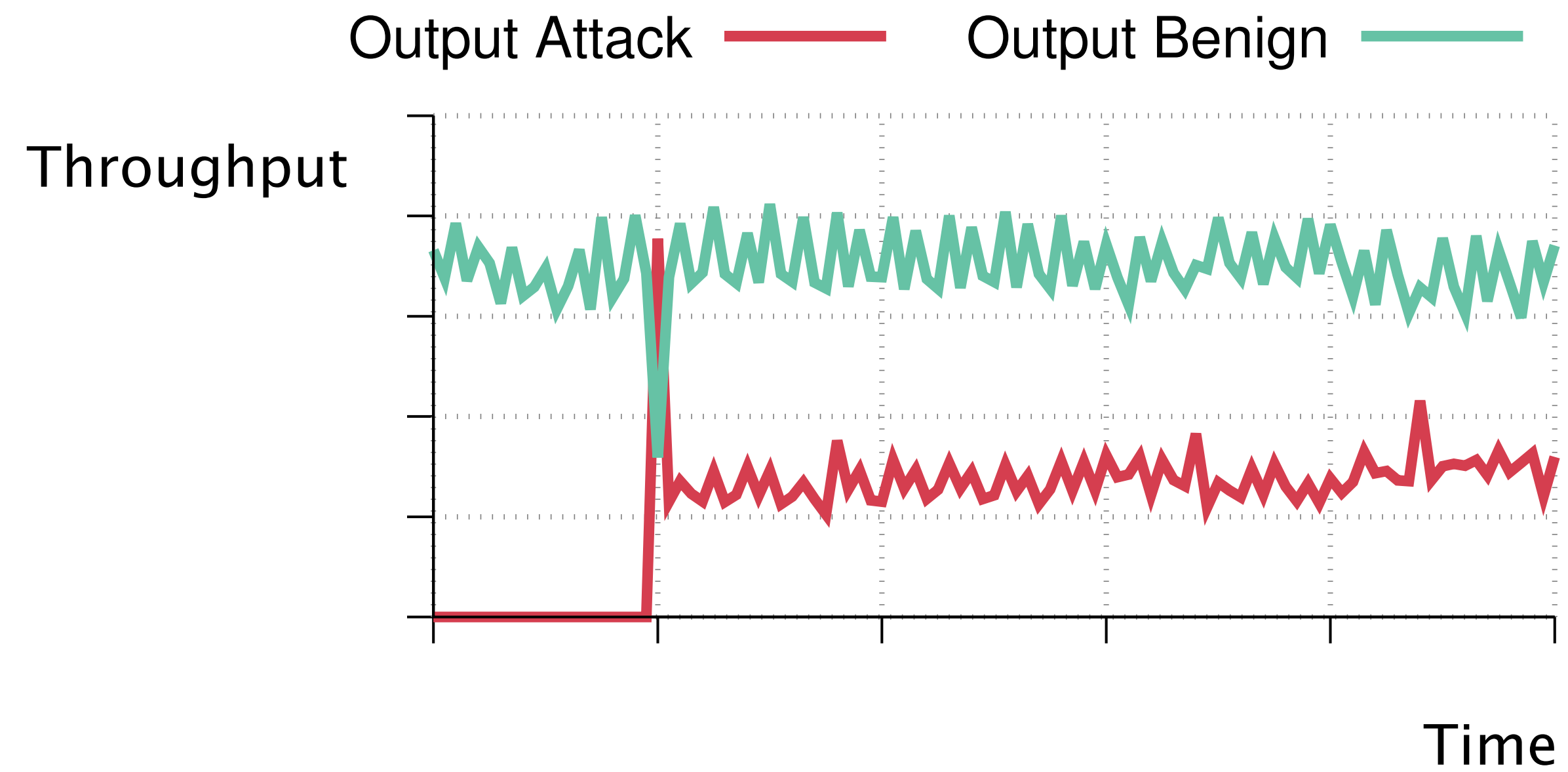
ACC-Turbo outperforms existing defenses and mitigates pulse-wave DDoS attacks

No defense



ACC-Turbo outperforms existing defenses and mitigates pulse-wave DDoS attacks

ACC-Turbo



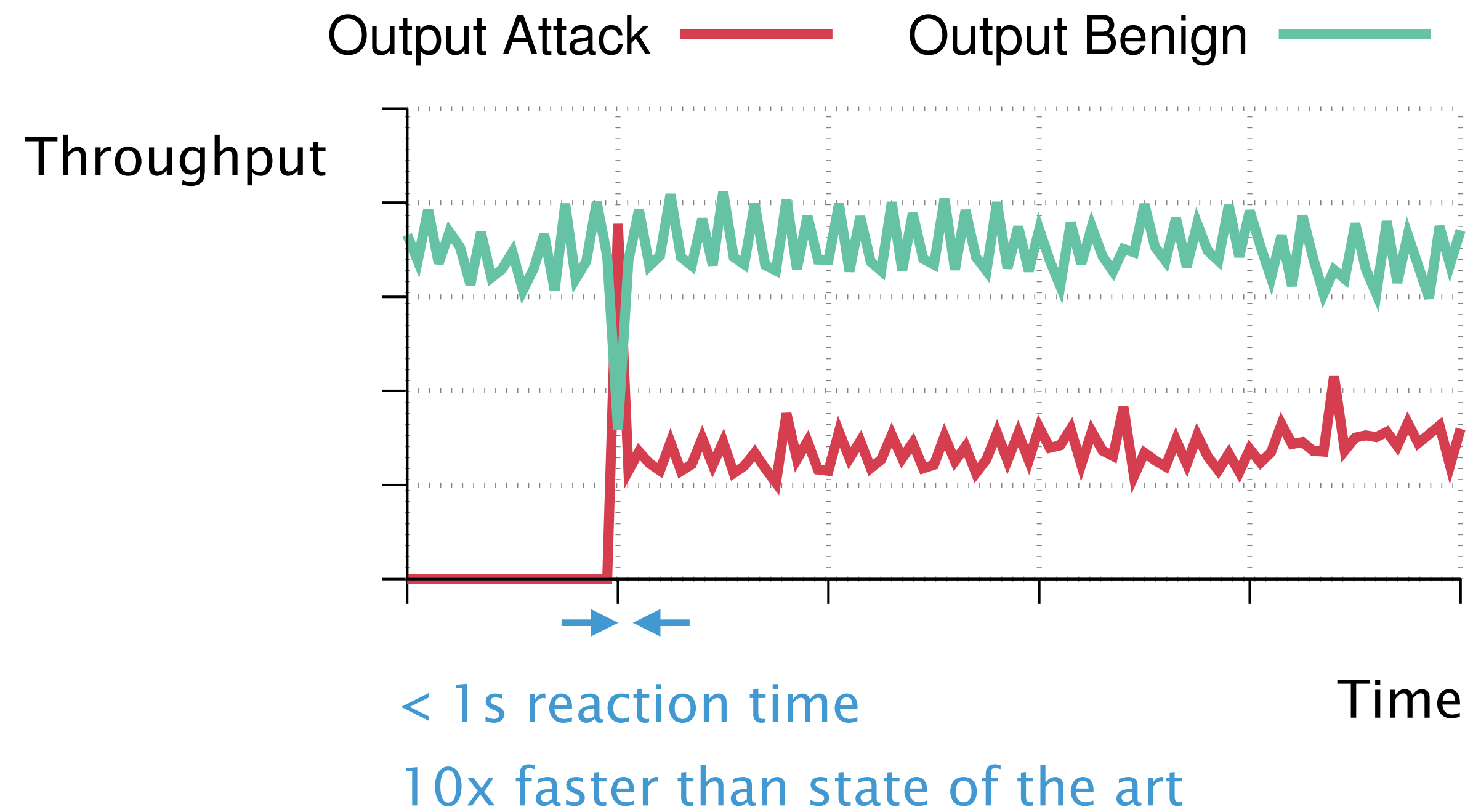
ACC-Turbo outperforms existing defenses and mitigates pulse-wave DDoS attacks

ACC-Turbo



ACC-Turbo outperforms existing defenses and mitigates pulse-wave DDoS attacks

ACC-Turbo



How to enable programmable scheduling on existing devices?

SP-PIFO

[NSDI'20]

Approximating
PIFO's scheduling

PACKS

[NSDI'25]

Incorporating
PIFO's admission

How to use it to improve the Internet's security?

ACC-Turbo

[SIGCOMM'22]

Mitigating
DDoS attacks

In-Network Congestion Management for Security and Performance

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How to use it to improve the Internet's security?

ACC-Turbo
[SIGCOMM'22]

Mitigating
DDoS attacks

Selected publications

NSDI '20

SP-PIFO: Approximating Push-In First-Out Behaviors using Strict-Priority Queues
A. Gran Alcoz, A. Dietmüller, L. Vanbever

SIGCOMM '22

Aggregate-Based Congestion Control for Pulse-Wave DDoS Defense
A. Gran Alcoz, M. Strohmeier, V. Lenders, L. Vanbever

HotNets '23

QVISOR: Virtualizing Packet Scheduling Policies
A. Gran Alcoz, L. Vanbever

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Principles for Internet Congestion Management
L. Brown, A. Gran Alcoz, F. Cangialosi, A. Narayan, M. Alizadeh, H. Balakrishnan,
E. Friedman, E. Katz-Bassett, A. Krishnamurthy, M. Schapira, S. Shenker

NSDI '25

Everything Matters in Programmable Packet Scheduling
A. Gran Alcoz, B. Vass, P. Namyar, B. Arzani, G. Rétvári, L. Vanbever

